Secrecy Performance of Wirelessly Powered Wiretap Channels


Published in:
IEEE Transactions on Communications

Document Version:
Early version, also known as pre-print

Queen's University Belfast - Research Portal:
Link to publication record in Queen's University Belfast Research Portal

General rights
Copyright for the publications made accessible via the Queen's University Belfast Research Portal is retained by the author(s) and / or other copyright owners and it is a condition of accessing these publications that users recognise and abide by the legal requirements associated with these rights.

Take down policy
The Research Portal is Queen's institutional repository that provides access to Queen's research output. Every effort has been made to ensure that content in the Research Portal does not infringe any person's rights, or applicable UK laws. If you discover content in the Research Portal that you believe breaches copyright or violates any law, please contact openaccess@qub.ac.uk.
Secrecy Performance of Wirelessly Powered Wiretap Channels

Xin Jiang, Student Member, IEEE, Caijun Zhong, Senior Member, IEEE, Xiaoming Chen, Senior Member, IEEE, Trung Q. Duong, Senior Member, IEEE, Theodoros Tsiftsis, Senior Member, IEEE, and Zhaoyang Zhang, Member, IEEE

Abstract—This paper considers a wirelessly powered wiretap channel, where an energy constrained multi-antenna information source, powered by a dedicated power beacon, communicates with a legitimate user in the presence of a passive eavesdropper. Based on a simple time-switching protocol where power transfer and information transmission are separated in time, we investigate two popular multi-antenna transmission schemes at the information source, namely maximum ratio transmission (MRT) and transmit antenna selection (TAS). Closed-form expressions are derived for the achievable secrecy outage probability and average secrecy rate for both schemes. In addition, simple approximations are obtained at the high signal-to-noise ratio (SNR) regime. Our results demonstrate that by exploiting the full knowledge of channel state information (CSI), we can achieve a better secrecy performance, e.g., with full CSI of the main channel, the system can achieve substantial secrecy diversity gain. On the other hand, without theCSI of the main channel, no diversity gain can be attained. Moreover, we show that the additional level of randomness induced by wireless power transfer does not affect the secrecy performance in the high SNR regime. Finally, our theoretical claims are validated by the numerical results.

Index Terms—Physical layer security, wireless power transfer, secrecy outage probability, average secrecy rate.

I. INTRODUCTION

Recently, the rapidly increasing demands for high data rate wireless services have put a tremendous pressure on the energy consumption of battery-powered mobile devices. Hence, how to prolong the lifetime of these energy-constrained mobile devices has become a critical problem to be addressed. Responding to this, energy harvesting techniques, which scavenge energy from ambient environment such as wind and solar have been proposed as a promising solution. Nevertheless, harvesting energy from nature resources depends heavily on the locations and weather conditions, which fails to generate stable energy output, hence may not be suitable to power wireless devices with strict quality of service requirements. As a result, a new energy harvesting paradigm, generally referred to as wireless power transfer (WPT), has gained considerable attentions. By exploiting the radio frequency (RF) signals as a means for energy transportation, WPT enables reliable and stable energy supplies to mobile devices. Since RF signals are widely used as a medium for information transmission, incorporating the feature of WPT into wireless communications networks has emerged as a hot topic, generally referred to as simultaneously wireless information and power transfer (SWIPT) systems, and significant research effects have been devoted to understand the fundamental performance limitation as well as design efficient SWIPT systems, see for instances [1]–[9] and references therein.

However, SWIPT systems are also vulnerable to potential security issues. This is because RF signals are shared by multiple nodes, which might be potential eavesdroppers. Recent research results show that compared to conventional cryptographic approaches, physical layer security is a better choice in energy and computation constrained systems, such as SWIPT systems [10]. The basic concept behind physical layer security is to exploit the physical layer characteristics of wireless channels to provide perfect secrecy. The work was pioneered by Wyner [11], which confirmed that perfect secrecy can be achieved when the quality of the wiretap channel is a degraded version of the main channel.

Recently, ensuing security in SWIPT systems have gained increasing attentions. In [12], [13], the authors presented the optimal beamforming design and power allocation scheme for multiple-input single-output (MISO) systems in the presence of passive eavesdroppers, later in [14], the issue of uncertain eavesdroppers was tackled, where joint optimization of information and energy beamforming and power allocation were studied. Latest works have considered the security issue in more sophisticated SWIPT systems, such as relay [15], multicast [16], cognitive radio [17] and OFDMA systems [18]. A common of these works is that they consider hybrid network architecture, where the information source also acts as the energy source. However, as analyzed in [19], the
harvested energy from hybrid networks is general infeasible to power larger devices such as smartphones, tablets and laptops. Responding to this, a novel network architecture was proposed in [20], where a dedicated station called power beacon (PB) is incorporated into the wireless network to power mobile devices. Very recently, the secrecy performance of device-to-device (D2D) communications in energy harvesting cognitive cellular networks has been investigated in [21], where the D2D transmitter first harvests energy from PBs, then performs secure transmission to the desired D2D receiver. Thus far, secure communications in SWIPT systems with dedicated PB remain largely an uncharted area. Motivated by this, we consider a point-to-point four-node wirelessly powered wiretap channel consisting of a dedicated PB, an energy constrained information source and a legitimate user in the presence of a passive eavesdropper. It is assumed that the source has no transmit power of its own, hence entirely relies on the external energy charging via wireless power transfer from the PB. For such systems, we present a comprehensive analysis on the achievable secrecy performance. It is worth pointing out that, unlike in the conventional communications systems, where the transmit power is constant, the use of WPT effectively makes the available source transmit power a random variable. In addition, since the transmit power affects both the signals observed at the legitimate user and eavesdropper, the effective signal-to-noise ratios (SNRs) of the main and wiretap channels become correlated, making the secrecy performance analysis much more challenging.

The main contributions of this work can be summarized as follows:

- For enhancing wireless security, we propose simple diversity transmission schemes at the information source, namely, maximum ratio transmission (MRT) and transmit antenna selection (TAS). In particular, for the TAS scheme, depending on the required channel state information (CSI), three different selection criteria are devised.
- For all schemes, closed-form expressions for secrecy outage probability are derived, which enable efficient evaluation of the achievable secrecy performance. Furthermore, simple and informative high SNR approximations are presented. The analytical results suggest that the achievable secrecy performance depends heavily on the available CSI at the source. With the CSI of the main channel, the system attains full secrecy diversity gain, while only unit secrecy diversity order can be achieved with only the CSI of the wiretap channel. In addition, the best performance is achieved when both the CSI of main channel and wiretap channel are available.
- For all schemes, closed-form expressions for average secrecy rate and high SNR approximations are also derived. Our results indicate that, all the schemes attain the same high SNR slope of one and distinct high SNR power offset. Moreover, increasing the number of transmit antennas improves the secrecy rate. However, the gain gradually diminishes when the number of transmit antennas is not large.
- Based on the simple high SNR expressions, the optimal time switching ratio \( \theta \) is studied. It was shown that there exists a unique \( \theta \) maximizing the secrecy throughput.

For the special single-antenna source case, closed-form expressions for the optimal \( \theta \) are obtained.
- We show that, the randomness of source transmit power induced by WPT does not affect the secrecy diversity order and the high SNR slope.

The remainder of the paper is organized as follows. Section II introduces the system model and proposes several transmission schemes. Section III provides an analytical study on the achievable secrecy outage probability of the proposed schemes, while Section IV investigates the average secrecy rate of the system. Numerical results and discussions are presented in Section V. Finally, Section VI concludes the paper and summarizes the key findings.

Notation: We use bold lower case letters to denote vectors and lower case letters to denote scalars; \( \| h \| \) denotes the Frobenius norm; \( E \{ x \} \) stands for the expectation of the random variable \( x \) and \( \lfloor x \rfloor \) denotes \( \max(0, x) \); \( T \) denotes the transpose operator and \( \dagger \) denotes the conjugate operator. \( I_k \) is the identity matrix of size \( k \). \( \Gamma(x) \) is the gamma function \([22, Eq. (8.31)]\), \( \Gamma(\alpha, x) \) is the upper incomplete gamma function \([22, Eq. (8.350.2)]\), and \( \gamma(\alpha, x) \) is the lower incomplete gamma function \([22, Eq. (8.350.1)]\). \( \psi(x) \) denotes the Euler psi function \([22, Eq. (8.36)]\) and \( K_n(x) \) is the \( n \)-th order modified Bessel function of the second kind \([22, Eq. (8.407.1)]\). \( S_{\alpha, \beta}(x) \) denotes the Lambert function \([22, Eq. (8.570.2)]\) and \( G_{\mu, \nu}^\alpha(x) \) denotes the Meijer G-function \([22, Eq. (9.301)]\).

II. SYSTEM MODEL

We consider a four-node wirelessly powered wiretap channel consisting of one PB, one information source Alice and one legitimate user Bob in the presence of one eavesdropper Eve as shown in Fig. 1. We assume that the source is equipped with \( N \) antennas, while the other three nodes are equipped with a single antenna.\(^1\) Quasi-static fading is assumed, such that the channel coefficients remain unchanged during each transmission block but vary independently between different blocks.

We adopt the time-sharing protocol proposed in [2]. Hence, a complete transmission slot with time duration of \( T \) is divided into two orthogonal sub-slots, i.e., the first one for power transfer with time duration of \( \theta T \) with \( \theta (0 < \theta < 1) \) being the time switching ratio, and the second one for information transmission with time duration of \( (1 - \theta)T \).

During the first phase, the PB sends an energy signal to Alice, and the received energy signal at Alice \( y_s \) can be expressed as

\[
y_s = \sqrt{P_S} h_p x_s + n_s,
\]

where \( P_S \) denotes the transmit power of the PB, \( x_s \) is the energy signal with unit power, \( n_s \) is an \( N \)-dimensional additive

\(^1\)The considered multi-antenna source model is relevant to the scenarios where an energy constrained multi-antenna sensor node \([7]\) transmits confidential information to a single-antenna fusion center or an energy constrained multi-antenna transmitter \([24]\) performs secrecy transmission to a single-antenna receiver.
white Gaussian noise (AWGN) vector with $E\{n_n\} = N_0 I$. The $N \times 1$ vector $h_p$ denotes the power transfer channel from PB to Alice. Due to relatively short distance between the two terminals, the line-of-sight propagation exists. Hence, the Nakagami-$m$ distribution is used to model the power transfer channel, i.e., the amplitude of each element of $h_p$ follows Nakagami-$m$ distribution with shape parameter $m$ and average power $\lambda_p$.\footnote{In the presence of line-of-sight effect, Rician fading is commonly used in literature. However, the analysis with Rician fading is much more involved. As such, we shall use Nakagami-$m$ fading model, since the Nakagami-$m$ distribution provides very accurate approximation to the Rician distribution.}

Therefore, at the end of the first phase, the total harvested energy within duration $\theta T$ can be expressed as

$$E = \eta P_S \|h_p\|^2 \theta T,$$

where $\eta (0 < \eta < 1)$ denotes the energy conversion efficiency.

Since the source communicates with the legitimate user during the second phase with duration $(1-\theta)T$, the transmit power can be computed as

$$P = \frac{E}{(1-\theta)T} = \eta P_S \|h_p\|^2 \frac{\theta}{1-\theta}.$$

To exploit the benefits of multiple antennas at Alice, different transmission schemes can be adopted. In this work, we consider two popular transmission schemes, namely MRT and TAS. The implementation of MRT and TAS requires different types of CSI. For MRT, only the CSI of the main channel is required. While for TAS, partial CSI of the main channel or the wiretap channel is required. In practice, the CSI of the main channel can be estimated at Bob, and then feedback to Alice. On the other hand, the CSI of the wiretap channel can be obtained when the eavesdropper is active in the network, a scenario that is particularly applicable in the networks combining multicast and unicast transmissions where the terminals play dual roles as legitimate receivers for some signals and eavesdroppers for others [25], [26].

A. Maximum Ratio Transmission (MRT)

For the MRT scheme, Alice aims at maximizing the reception quality of the main channel by making use of a channel.

$\eta_1, \eta_2 < 1$. Hence, the Nakagami-$m$ distribution provides very accurate approximation to the Rician distribution.

match beam, as such, the received signal $y_M$ at Bob can be written as

$$y_M = \sqrt{\eta} h_M^T w x + n_M,$$

where $x$ denotes the information symbol with unit energy; $N \times 1$ vector $h_M$ denotes the main channel from Alice to Bob, whose elements are circularly symmetric complex Gaussian random variables (RVs) with zero mean and variance $\lambda_M$; $n_M$ denotes the AWGN with zero mean and variance $N_0$; $w$ is the MRT vector given by $w = \frac{h_M^*}{\|h_M\|^2}$. Similarly, the received signal at Eve $y_W$ can be expressed as

$$y_W = \sqrt{\eta} h_W^T w x + n_W,$$

where the $N \times 1$ vector $h_W$ denotes the wiretap channel from Alice to Eve, whose elements are circularly symmetric complex Gaussian RVs with zero mean and variance $\lambda_W$, and $n_W$ denotes the AWGN with zero mean and variance $N_0$. As such, the instantaneous SNR at Bob $\gamma_M$ and at Eve $\gamma_W$ are given by

$$\gamma_M = \eta P_S \|h_M\|^2 \|h_M\|^2 \frac{\theta}{1-\theta},$$

and

$$\gamma_W = \eta P_S \|h_W\|^2 \|h_W\|^2 \frac{\theta}{1-\theta}.$$

respectively.

B. Transmit Antenna Selection (TAS)

TAS is another low-complexity transmission scheme. In this work, we consider three different selection criteria as elaborated below.

1) Criterion 1: In this case, the antenna with the maximum gain of main channel is selected, i.e.,

$$k = \arg \max_{i=1,\ldots,N} |h_{id}|^2,$$

where $h_{id}$ is the $i$-th element of main channel $h_M$. It is worth noting that best antenna selection according to the above criterion implies a random antenna selection for the wiretap channel because the main channel is independent of the wiretap channel.

2) Criterion 2: Instead of maximizing the gain of main channel, we now intend to minimize the gain of wiretap channel. As such, the best antenna is selected according to the following criterion:

$$k = \arg \min_{i=1,\ldots,N} |h_{ie}|^2,$$

where $h_{ie}$ is the $i$-th element of the wiretap channel $h_W$.

3) Criterion 3: Since the secrecy performance of system depends on the quality of both the main channel and wiretap channel, we now propose the third selection criterion which picks the antenna maximizing the ratio of main channel gain and wiretap channel gain, i.e.,

$$k = \arg \max_{i=1,\ldots,N} \left( \frac{|h_{id}|^2}{|h_{ie}|^2} \right).$$
The exact secrecy outage probability of the TAS scheme can be expressed as:

\[
P_{\text{out}}^{\text{TAS}}(R_S) = P(C_S < R_S).
\]

A. MRT

We start with the MRT scheme, and we have the following key result:

**Theorem 1:** The exact secrecy outage probability of the MRT scheme can be expressed in closed-form as

\[
P_{\text{out}}^{\text{MRT}}(R_S) = 1 - \frac{2}{\Gamma(mN)} \sum_{k=0}^{N-1} \frac{\lambda_M \lambda_W}{\lambda_M + k_2 \lambda_W(k+1)} \times \left( \frac{m(k+1)(k_2-1)}{k_1 \lambda_M \lambda_P} \right)^N K_{mN} \left( \frac{m(k+1)(k_2-1)}{k_1 \lambda_M \lambda_P} \right).
\]

where \( k_1 = \frac{N_{\text{PB}}}{N_0} \lambda_M^{-1} \) and \( k_2 = 2R_S \).

Proof: See Appendix A.

**Theorem 1** presents an exact closed-form expression for the secrecy outage probability, which can be efficiently evaluated. However, the expression is too complicated to yield any insights. Motivated by this, we now look into the asymptotic regime, where simple expressions can be obtained.

For the asymptotic high SNR regime, we assume that \( \lambda_M \to \infty \) with an arbitrary \( \lambda_W \). Such a scenario has been widely adopted in the literature, see for instance [27]–[30]. In practice, this occurs when the quality of the main channel is much better than wiretap channel, i.e., Bob is relatively close to Alice while Eve is far away from Alice or the wiretap channel undergoes severe small-scale and large-scale fading effects. In the following, we characterize the two key performance parameters governing the secrecy outage probability in the high SNR regime, i.e., secrecy diversity order \( G_d \) and secrecy array gain \( G_a \) defined by [31]

\[
P_{\text{out}}^{\text{MRT}}(R_S) = (G_a \lambda_M)^{-G_d}.
\]

**Proposition 1:** In the high SNR regime, i.e., \( \lambda_M \to \infty \), the secrecy outage probability of the MRT scheme can be approximated by

\[
P_{\text{out}}^{\infty}(R_S) = \sum_{k=0}^{N} \frac{\Gamma(mN - k)}{\Gamma(mN)} \left( \frac{m(k_2 - 1)}{k_1 \lambda_W \lambda_P} \right)^k \left( \frac{k_2 \lambda_W}{\lambda_M} \right)^N.
\]

Proof: See Appendix B.

It is evident from (17) that the system achieves a secrecy diversity order of \( N \). In addition, we observe the intuitive effect of the position of nodes on the secrecy outage probability. For instance, the secrecy outage probability decreases when the PB is close to the source, i.e., large \( \lambda_P \). It is also easy to see that the high SNR secrecy outage probability \( P_{\text{out}}^{\infty}(R_S) \) is a decreasing function with respect to \( \frac{1}{\lambda_M} \), indicating that increasing the transmit power of the PB is always beneficial.

B. TAS Criterion 1

We now move to the TAS Criterion 1 scheme, and we obtain the following key result:

**Theorem 2:** The exact secrecy outage probability of TAS Criterion 1 scheme can be expressed in closed-form as

\[
P_{\text{out}}^{\text{TAS1}}(R_S) = 1 - \frac{2}{\Gamma(mN)} \sum_{k=0}^{N-1} \frac{\lambda_M \lambda_W}{\lambda_M + k_2 \lambda_W(k+1)} \times \left( \frac{m(k+1)(k_2-1)}{k_1 \lambda_M \lambda_P} \right)^N K_{mN} \left( \frac{m(k+1)(k_2-1)}{k_1 \lambda_M \lambda_P} \right).
\]

**Proposition 2:** In the high SNR regime, i.e., \( \lambda_M \to \infty \), and secrecy diversity order \( G_d \), the secrecy outage probability of the TAS scheme can be approximated by

\[
P_{\text{out}}^{\infty}(R_S) = \sum_{k=0}^{N} \frac{\Gamma(mN - k)}{\Gamma(mN)} \left( \frac{m(k_2 - 1)}{k_1 \lambda_W \lambda_P} \right)^k \left( \frac{k_2 \lambda_W}{\lambda_M} \right)^N.
\]

Although the energy transfer efficiency of state of the art technique is low, it is still of both theoretical and practical interests in wirelessly powered communications systems due to the following reasons: First, the energy transfer efficiency can be significantly improved by adopting multiple antenna technology and the PB assisted WPC architecture. Second, the effective SNR could be still reasonably high even if the energy transfer efficiency is low. Third, the key insights obtained from high SNR analysis provide useful guidance for practical system design.

---

*IEEE TRANSACTIONS ON COMMUNICATIONS, VOL. XX, NO. XX, XX 2016 4*

---

\( R \) can be expressed mathematically as

\[
\gamma_M \quad \text{and at Eve} \quad \gamma_W
\]

are characterized in terms of the secrecy outage probability, which can be efficiently evaluated. However, the expression is too complicated to yield any insights. Motivated by this, we now look into the asymptotic regime, where simple expressions can be obtained.

For the asymptotic high SNR regime, we assume that \( \lambda_M \to \infty \) with an arbitrary \( \lambda_W \). Such a scenario has been widely adopted in the literature, see for instance [27]–[30]. In practice, this occurs when the quality of the main channel is much better than wiretap channel, i.e., Bob is relatively close to Alice while Eve is far away from Alice or the wiretap channel undergoes severe small-scale and large-scale fading effects. In the following, we characterize the two key performance parameters governing the secrecy outage probability in the high SNR regime, i.e., secrecy diversity order \( G_d \) and secrecy array gain \( G_a \) defined by [31]

\[
P_{\text{out}}^{\infty}(R_S) = (G_a \lambda_M)^{-G_d}.
\]

**Proposition 1:** In the high SNR regime, i.e., \( \lambda_M \to \infty \), the secrecy outage probability of the MRT scheme can be approximated by

\[
P_{\text{out}}^{\infty}(R_S) = \sum_{k=0}^{N} \frac{\Gamma(mN - k)}{\Gamma(mN)} \left( \frac{m(k_2 - 1)}{k_1 \lambda_W \lambda_P} \right)^k \left( \frac{k_2 \lambda_W}{\lambda_M} \right)^N.
\]

Proof: See Appendix B.

It is evident from (17) that the system achieves a secrecy diversity order of \( N \). In addition, we observe the intuitive effect of the position of nodes on the secrecy outage probability. For instance, the secrecy outage probability decreases when the PB is close to the source, i.e., large \( \lambda_P \). It is also easy to see that the high SNR secrecy outage probability \( P_{\text{out}}^{\infty}(R_S) \) is a decreasing function with respect to \( \frac{1}{\lambda_M} \), indicating that increasing the transmit power of the PB is always beneficial.
This article has been accepted for publication in a future issue of this journal, but has not been fully edited. Content may change prior to final publication. Citation information: DOI 10.1109/TCOMM.2016.2592529, IEEE Transactions on Communications

Criterion 2 scheme can be expressed in closed-form as that of Proposition 1 yields the desired result.

The main and wiretap channel capacities, and the transmit secrecy performance is determined by the difference between outage probability. This is because that in wiretap channels, the power does not affect the scaling behavior of the secrecy current work, we notice that the randomness of the transmit see that

\[ P_{\text{TAS}}(R_S) = \left( \frac{1}{N + \frac{1}{mN - 1}} \right) \left( \frac{m(k_2 - 1)}{k_1k_2\lambda_W\lambda_P} \right) . \] (21)

\( \lambda_M \rightarrow \infty \)

D. TAS Criterion 3

We now analyze the secrecy outage probability of the system with the TAS Criterion 3 scheme.

Theorem 4: The secrecy outage probability of TAS Criterion 3 scheme can be approximated by

\[ P_{\text{TAS}}(R_S) \approx \left( \frac{k_2}{k_2 + \frac{1}{\lambda_M}} \right)^N . \] (22)

Proof: See Appendix F.

Having obtained the outage probability of TAS Criterion 3 scheme, we now look into the asymptotic regime.

Proposition 4: In the high SNR regime, i.e., \( \lambda_M \rightarrow \infty \), the secrecy outage probability of TAS Criterion 3 scheme can be approximated as

\[ P_{\text{TAS}}(R_S) = \left( \frac{k_2\lambda_W}{\lambda_M} \right)^N . \] (23)

Proof: The proof is straightforward, hence omitted.

As expected, the system achieves a secrecy diversity order of \( N \). Recall the high SNR outage probability of the MRT scheme, and noticing that

\[ 1 + \sum_{k=1}^{mN-k} \left( \frac{m(k_2-1)}{k_1k_2\lambda_W\lambda_P} \right)^k > 1 \], we observe that the TAS Criterion 3 scheme outperforms the MRT scheme. This is reasonable since the TAS Criterion 3 scheme considers both the CSI of \( \mathbf{h}_M \) and \( \mathbf{h}_W \), while only the CSI of the \( \mathbf{h}_M \) is utilized in the MRT scheme.

E. Optimization of the Time Switching Ratio \( \theta \)

From the analytical expressions derived in previous subsections, we are ready to study the optimization of time switching ratio \( \theta \). Specifically, we adopt the effective secrecy throughput as the performance measure as in [34]. Hence, when the source transmits at a constant rate \( R_S \), the average secrecy throughput can be evaluated by

\[ R = (1 - P_{\text{out}})R_S(1 - \theta). \] Also, due to space limitation and for illustrative purpose, we only focus on the MRT scheme, while other cases will be numerically illustrated in Section V.

Using the high SNR approximation given in (17), the effective secrecy throughput of the MRT scheme can be expressed
as
\[ \tau(\theta) = (1 - P_{MRT}^{\infty}(\theta))R_S \frac{(1 - \theta)T}{T} = (1 - P_{MRT}^{\infty}(\theta))R_S (1 - \theta). \] (24)

Hence, the optimal \( \theta^* \) is the solution of the following optimization problem:
\[ \theta^* = \arg \max_{\theta} \tau(\theta) \quad \text{s.t.} \quad 0 < \theta < 1. \] (25)

To this end, we have the following key result:

**Proposition 5:** Consider a polynomial
\[ \sum_{k=0}^{N} a_k \left( \frac{1 - \theta}{\theta} \right)^k b(1 - \theta) = 0, \] (26)

where \( a_k = \frac{1}{\prod_{i=0}^{N} \frac{m(k-1)N_i}{\psi_i} \frac{m(k-1)}{mN}} \). Then, the optimal \( \theta^* \) is the unique root of the polynomial in (0,1).

**Proof:** Substituting (17) into (24), we obtain
\[ \tau(\theta) = \left( 1 - \sum_{k=0}^{N} a_k \left( \frac{1 - \theta}{\theta} \right)^k \right) b(1 - \theta), \] (27)

where \( b = R_S \). Thus, the derivative of \( \tau(\theta) \) with respect to \( \theta \) can be expressed as
\[ \frac{d\tau(\theta)}{d\theta} = \sum_{k=0}^{N} a_k b \left( \frac{1 - \theta}{\theta} \right)^k \left( 1 + \frac{k}{\theta} \right) - b. \] (28)

It is easy to show that \( \frac{d\tau(\theta)}{d\theta} \) is a monotonically decreasing function with respect to \( \theta \). When \( \theta \) approaches \( 0 \), \( \frac{d\tau(\theta)}{d\theta} \) approaches \( +\infty \) and when \( \theta \) approaches \( 1 \), \( \frac{d\tau(\theta)}{d\theta} \) approaches \( 0 \). Therefore, there exists a unique \( \theta^* \in (0,1) \) with \( \frac{d\tau(\theta)}{d\theta} = 0 \), where \( \tau(\theta) \) attains its maximum value.

In general, due to the complexity of the involved expression, deriving a closed-form solution for \( \theta^* \) is very challenging. However, for small \( N \), closed-form expressions for \( \theta^* \) can be obtained.

**Remark 1:** When \( N = 1 \), the optimal \( \theta^* \) is given by
\[ \theta^* = \sqrt{\frac{a_1}{1 + a_1 - a_0}}. \] (29)

**Remark 2:** When \( N = 2 \), the optimal \( \theta^* \) is given by
\[ \theta^* = \sqrt{\frac{q}{2}} + \sqrt{\left( \frac{q}{2} \right)^2 + \left( \frac{p}{3} \right)^2} + q - \sqrt{\left( \frac{q}{2} \right)^2 + \left( \frac{p}{3} \right)^2}, \] (30)

where \( p = \frac{3a_2 - a_1}{1 - a_0 + a_1 - a_2} \) and \( q = -\frac{2a_2}{1 - a_0 + a_1 - a_2} \).

**F. Comparison of the Proposed Protocols**

We now present a more detailed performance comparison for the proposed schemes at the high SNR regime as summarized in Table I on the top of the next page. In general, the secrecy performance depends heavily on the available CSI at the source. The more CSI available, the better the secrecy performance. However, in terms of high SNR protocols investigated shown in Table I. We list the CSI that is required for each protocol and investigate the secrecy diversity order and the outage performance according to the asymptotic secrecy outage probability.

**IV. AVERAGE SECRECY RATE**

In this section, we focus on the average secrecy rate performance of the system. For both transmission schemes, new closed-form expressions for the exact and asymptotic average secrecy rate are presented. Based on which, the impacts of multiple antennas on the secrecy performance are characterized in terms of the high SNR slope and the high SNR power offset.

Starting from the definition, the average secrecy rate can be expressed as
\[ C = E[(\log_2(1 + \gamma^*_M) - \log_2(1 + \gamma^*_W))]\] (31)

We now study the achievable secrecy rate of different transmission schemes in the following.

**A. MRT**

**Theorem 5:** The exact average secrecy rate of MRT scheme can be expressed in closed-form as (32) on the top of the next page.

**Proof:** See Appendix G.

While (32) provides the exact average secrecy rate, the expression is too complex to yield insightful information. Motivated by this, we now look into the high SNR regime, where the secrecy rate is dictated by two key parameters known as the high SNR slope \( S_{\infty} \) and the high SNR power offset \( L_{\infty} \) [31], i.e.,
\[ C_{\infty} = S_{\infty} (\log_2(\lambda_M) - L_{\infty}). \] (33)

To this end, we have the following key result:

**Proposition 6:** For the MRT scheme, the high SNR slope \( S_{\infty} \) is given by
\[ S_{\infty}^{MRT} = 1, \] (34)

and the high SNR power offset \( L_{\infty}^{MRT} \) is given by
\[ L_{\infty}^{MRT} = -\frac{1}{\ln 2} \left( \ln \frac{k_1 \lambda_P}{m} + \psi(N) + \psi(mN) \right) + \frac{1}{\Gamma(mN) \ln 2} G_{13}^{31} \left( \frac{m}{k_1 \lambda_W} \left| \frac{0}{mN, N, 0} \right. \right). \] (35)

**Proof:** See Appendix H.

We note that the high SNR slope is also known as the maximum multiplexing gain or the number of degrees of freedom [35]. According to (34), the high SNR slope is one because we assume that the legitimate user only employs a single antenna in this system. In addition, we observe that the wiretap channel only reflects the high SNR power offset. It is also easy to see that the asymptotic average secrecy rate is an increasing function with respect to \( k_1 \) and \( \lambda_P \), indicating that increasing the transmit power of the PB is always beneficial and the secrecy rate increases when PB is close to the source.
TABLE I: Comparison of the proposed schemes

<table>
<thead>
<tr>
<th>Scheme</th>
<th>CSI requirement</th>
<th>Secrecy diversity order</th>
<th>Outage performance</th>
</tr>
</thead>
<tbody>
<tr>
<td>MRT</td>
<td>$H_M$</td>
<td>$N$</td>
<td>Second best</td>
</tr>
<tr>
<td>TAS Criterion 1</td>
<td>Index of the entry of $H_M$</td>
<td>$N$</td>
<td>Third best</td>
</tr>
<tr>
<td>TAS Criterion 2</td>
<td>Index of the entry of $H_W$</td>
<td>1</td>
<td>Worst</td>
</tr>
<tr>
<td>TAS Criterion 3</td>
<td>$H_M$ and $H_W$</td>
<td>$N$</td>
<td>Best</td>
</tr>
</tbody>
</table>

$$C_{\text{MRT}} = \frac{1}{\Gamma(mN)\ln 2} \left[ \sum_{k=0}^{N-1} \left( \frac{\lambda_P}{m} \right)^k + \frac{1}{\Gamma(mN-1)\ln 2} \right]$$

\[ \begin{align*}
C_{\text{TAS}} = & \frac{1}{\Gamma(mN)\ln 2} \left[ G_{13}^{31} \left( \frac{m}{k_1\lambda_P \lambda_M} \right) \left( mN, 0, 0 \right) \right] \\
& \quad - G_{13}^{31} \left( \frac{m}{k_1\lambda_P + \frac{N}{\lambda_W}} \right) \left( mN, 0, 0 \right) .
\end{align*} \]

B. TAS Criterion 1

We now consider the TAS Criterion 1 scheme, and we have the following key result:

**Theorem 6:** The exact average secrecy rate of TAS criterion 1 scheme can be expressed in closed-form as (36) on the top of the next page.

**Proof:** The proof follows similar lines as that of Theorem 5, hence is omitted.

We now look into the high SNR regime, and present the high SNR metrics in the following proposition.

**Proposition 7:** For the TAS criterion 1 scheme, the high SNR slope $S_{\infty}^{\text{TAS1}}$ is given by

$$S_{\infty}^{\text{TAS1}} = 1,$$  \hspace{1cm} (37)

and the high SNR power offset $L_{\infty}^{\text{TAS1}}$ is given by

$$L_{\infty}^{\text{TAS1}} = -\frac{1}{\ln 2} \left( \frac{\lambda_P}{m} + \psi(1) + \psi(mN) + \frac{1}{\Gamma(mN)\ln 2} G_{13}^{31} \left( \frac{m}{k_1\lambda_P \lambda_M} \right) \left( mN, 0, 0 \right) \right).$$  \hspace{1cm} (38)

**Proof:** The proof follows similar lines as that of Proposition 6, hence is omitted.

C. TAS Criterion 2

We now analyze the average secrecy rate of the system with the TAS Criterion 2 scheme, and we have the following key result:

**Theorem 7:** The exact average secrecy rate of TAS criterion 2 scheme can be expressed in closed-form as

$$\tilde{C}_{\text{TAS}2} = \frac{1}{\Gamma(mN)\ln 2} \left[ G_{13}^{31} \left( \frac{m}{k_1\lambda_P \lambda_M} \right) \left( mN, 0, 0 \right) \right]$$

\[ \begin{align*}
& \quad - G_{13}^{31} \left( \frac{m}{k_1\lambda_P + \frac{N}{\lambda_W}} \right) \left( mN, 0, 0 \right) .
\end{align*} \]

**Proof:** The proof follows similar lines as that of Theorem 5, hence is omitted.

Having obtained the exact average secrecy rate of TAS Criterion 2 scheme, we now look into the asymptotic regime.

**Proposition 8:** For the TAS criterion 2 scheme, the high SNR slope $S_{\infty}^{\text{TAS2}}$ is given by

$$S_{\infty}^{\text{TAS2}} = 1,$$  \hspace{1cm} (40)

and the high SNR power offset $L_{\infty}^{\text{TAS2}}$ is given by

$$L_{\infty}^{\text{TAS2}} = -\frac{1}{\ln 2} \left( \frac{\lambda_P}{m} + \psi(1) + \psi(mN) + \frac{1}{\Gamma(mN)\ln 2} G_{13}^{31} \left( \frac{m}{k_1\lambda_P \lambda_M} \right) \left( mN, 0, 0 \right) \right).$$  \hspace{1cm} (41)

**Proof:** The proof follows similar lines as that of Proposition 6, hence is omitted.

D. TAS Criterion 3

We now move to the TAS Criterion 3 scheme, and we obtain the following key result:

**Theorem 8:** The average secrecy rate of TAS criterion 3 scheme can be approximated by

$$\tilde{C}_{\text{TAS}3} \approx \frac{1}{\ln 2} \left( \frac{\lambda_M}{\lambda_P} \sum_{k=0}^{N-1} \frac{1}{\Gamma(k+1)\Gamma(N-k)} \times \right.$$

$$\left. \quad G_{13}^{33} \left( 1 + \frac{\lambda_M}{\lambda_W} \right) \left( -N, 1+k-N, 1+k-N \mid 0, 1+k-N, k-N \right) \right).$$  \hspace{1cm} (42)

**Proof:** See Appendix I.

We now look into the high SNR regime, and present the high SNR metrics in the following proposition.

**Proposition 9:** For the TAS criterion 3 scheme, the high SNR slope $S_{\infty}^{\text{TAS3}}$ is given by

$$S_{\infty}^{\text{TAS3}} = 1,$$  \hspace{1cm} (43)

and the high SNR power offset $L_{\infty}^{\text{TAS3}}$ is given by

$$L_{\infty}^{\text{TAS3}} = -\frac{1}{\ln 2} \left( \ln \lambda_W + \frac{1}{N} + \sum_{k=2}^{N} \left( -1 \right)^k \frac{N!}{k!} \right).$$  \hspace{1cm} (44)

**Proof:** See Appendix J.
\[ \hat{C}_{\text{TAS1}} = \frac{1}{1/(mN) \ln 2} \sum_{k=1}^{N} (-1)^{k+1} \binom{N}{k} \left[ G_{13}^\text{TAS1} \left( \frac{m}{k_1 \lambda_P} \frac{k}{\lambda_M} \right) \mid 0 \right] mN, 0, 0 \right] - \frac{1}{1/(mN) \ln 2} \sum_{k=1}^{N} (-1)^{k+1} \binom{N}{k} \left[ G_{13}^\text{TAS1} \left( \frac{m}{k_1 \lambda_P} \frac{k}{\lambda_M} + \frac{1}{\lambda_W} \right) \mid 0 \right] mN, 0, 0 \right]. \]  

(36)

V. NUMERICAL RESULTS

In this section, we present numerical results to verify the theoretical expressions. Unless otherwise stated, we set the source transmission rate as \( R_S = 1 \text{ bit/s/Hz} \), the energy conversion efficiency as \( \eta = 0.8 \) and the time switching ratio as \( \theta = 0.5 \). The Nakagami-\( m \) parameter is set to be \( m = 4 \), which corresponds to a Rician factor of \( K = 3 + \sqrt{12} \). The transmit power of the PB to the noise ratio as \( \frac{P_0}{\lambda} = 10 \text{ dB} \), the channel variance as \( \lambda_P = 1 \) and \( \lambda_W = 10 \). Also, we set \( \rho = \frac{P_0}{N_0} \lambda_M \) and \( \rho_1 = \frac{P_0}{N_0} \lambda_W \) to denote the average SNR of the main channel and the wiretap channel, respectively.

Fig. 2: Secrecy outage probability versus \( \rho \) with different \( N \) for the MRT scheme.

Fig. 2 plots the secrecy outage probability versus \( \rho \) with different \( N \) for the MRT scheme. As illustrated, the analytical results are in exact agreement with the Monte Carlo simulations, which demonstrates the correctness of the analytical expression. In addition, the high SNR results accurately predict the secrecy diversity order and the secrecy array gain, and increasing \( N \) substantially enhances the secrecy outage performance by achieving a higher secrecy diversity gain.

Fig. 3: Secrecy outage probability versus \( \rho \) for different TAS schemes with \( N = 3 \).

Fig. 3 examines the average secrecy rate of the MRT scheme. We observe that the high SNR slope of all the curves is one, which corroborates the theoretical analysis presented in previous section. Nevertheless, increasing \( N \) improves the average secrecy rate by decreasing the high SNR power offset. However, the benefit of increasing the number of antennas \( N \) gradually diminishes when \( N \) is sufficiently large.

Fig. 4: Average secrecy rate versus \( \rho \) with different \( N \) for the MRT scheme.

Fig. 4 investigates the impact of distance between Alice and Eve on the average secrecy rate for different TAS schemes. As expected, the TAS Criterion 3 scheme always attains the best performance. At the low \( \rho_1 \) regime, namely, relatively large distance between Alice and Eve (or small \( \lambda_W \)), the TAS Criterion 1 scheme outperforms the TAS Criterion 2 scheme. However, the opposite holds when the distance decreases, i.e., \( \lambda_W \) becomes large. This phenomenon is quite intuitive, since when the average channel gain of the wiretap channel is better than that of the main channel, selecting the worst antenna as per the TAS Criterion 2 scheme substantially degrades the capacity of the wiretap channel, thereby resulting in a larger secrecy performance gain than the TAS Criterion 1 scheme.

Fig. 5: Average secrecy rate versus \( \rho_1 \) for different TAS schemes with \( N = 3 \).

Fig. 5 examines the average secrecy rate of the MRT scheme. We observe that the high SNR slope of all the curves is one, which corroborates the theoretical analysis presented in previous section. Nevertheless, increasing \( N \) improves the average secrecy rate by decreasing the high SNR power offset. However, the benefit of increasing the number of antennas \( N \) gradually diminishes when \( N \) is sufficiently large.

Fig. 6: Secrecy outage probability versus \( \rho \) with different \( \lambda_W \) for the MRT scheme.

Fig. 6 investigates the impact of distance between Alice and Eve on the average secrecy rate for different TAS schemes. As expected, the TAS Criterion 3 scheme always attains the best performance. At the low \( \rho_1 \) regime, namely, relatively large distance between Alice and Eve (or small \( \lambda_W \)), the TAS Criterion 1 scheme outperforms the TAS Criterion 2 scheme. However, the opposite holds when the distance decreases, i.e., \( \lambda_W \) becomes large. This phenomenon is quite intuitive, since when the average channel gain of the wiretap channel is better than that of the main channel, selecting the worst antenna as per the TAS Criterion 2 scheme substantially degrades the capacity of the wiretap channel, thereby resulting in a larger secrecy performance gain than the TAS Criterion 1 scheme.
Fig. 5: Average secrecy rate versus \( \rho_1 \) for different TAS schemes with \( N = 3 \).

Fig. 6: Average secrecy rate versus \( \rho \) for different TAS schemes.

Fig. 7: Effective throughput versus \( \theta \) for different schemes with \( N = 2 \) and \( \lambda_M = 1 \).

forms the TAS Criterion 1. The reason is that we have set \( \lambda_M = 1 \) and \( \lambda_W = 10 \) in the simulations, a scenario where the average gain of the main channel is worse than the wiretap channel. Hence, the main channel capacity enhancement due to selection of the strongest channel in TAS Criterion 1 scheme is rather insignificant compared to the wiretap channel capacity degradation due to the selection of the weakest channel in TAS Criterion 2 scheme.

VI. CONCLUSION

In this paper, we have investigated the secrecy performance of the wirelessly powered wiretap channels. For both MRT and TAS schemes, exact analytical expressions and asymptotic approximations are presented, which facilitate the extraction of key insights of the achievable secrecy outage probability and average secrecy rate performance. The findings of the paper suggest that, with CSI of the main channel (e.g., MRT and TAS Criteria 1 and 3), the system can achieve substantial secrecy diversity gain. On the other hand, without the CSI of the main channel (e.g., TAS Criterion 2), no diversity gain can be attained, which indicates the critical importance of CSI in the design of practical systems.
\[ \sum_{y} P \text{ distribution with shape parameter } \lambda \text{ and scale parameter } \frac{y}{\lambda} \]

We start by expressing the SNR given in (6) and (7) as
\[ k \frac{\sum_{y} P}{m} = k_{1} y_{w} h_{m} + \frac{y_{m}}{y_{w}} \]

where \( k_{1} = \frac{n_{P}}{N_{r}} \), \( y_{w} = ||h_{w}||^{2} \), \( y_{m} = ||h_{m}||^{2} \) and \( y_{w} = \frac{h_{w}^{\dagger} h_{w}}{||h_{w}||^{2}} \). It is straightforward to show that the probability density function (pdf) of \( y_{w} \) follows a gamma distribution with shape parameter \( mN \) and scale parameter \( \lambda P/m \) by [36]
\[ f_{y_{w}}(x) = \frac{1}{\Gamma(mN)} \left( \frac{m}{\lambda P} \right)^{mN} x^{mN-1} e^{-\frac{x}{\lambda P}} \] \[ (46) \]

and the pdf of \( y_{m} \) follows a chi-square distribution with \( 2N \) degrees of freedom given by [37]
\[ f_{y_{m}}(x) = \frac{x^{N-1}}{\lambda_{M}^{N} \Gamma(N)} e^{-\frac{x}{\lambda_{M}}} \] \[ (47) \]

In addition, according to [38], \( y_{w} \) follows an exponential distribution with pdf
\[ f_{y_{w}}(x) = \frac{1}{\lambda_{w}} e^{-\frac{x}{\lambda_{w}}} \] \[ (48) \]

and is independent of \( y_{m} \). As such, the secrecy outage probability can be written as
\[ P_{out}^{MRT}(R_{S}) = 1 - P \left( \frac{1 + k_{1} y_{w} y_{m}}{1 + k_{1} y_{w} y_{w}} \geq k_{2} \right) \] \[ (49) \]

where \( k_{2} = 2R_{e} \). Conditioned on \( y_{w} \) and \( y_{w} \), with the help of [22, Eq. (3.351.2)], we obtain
\[ P_{out}^{MRT}(R_{S}|y_{w}, y_{w}) = 1 - \int_{k_{2}}^{\infty} \frac{x^{N-1}}{\left( \sum_{k=0}^{N} \binom{k}{k} x_{1}^{x_{1}} x_{2}^{x_{2}} \right)} e^{-\frac{x_{1} + x_{2}}{\lambda_{w}}} dx \]
\[ (50) \]

By applying the binomial expansion \( (x_{1} + x_{2})^{m} = \sum_{k=0}^{m} \binom{m}{k} x_{1}^{k} x_{2}^{m-k} \), (50) can be further expressed as
\[ P_{out}^{MRT}(R_{S}|y_{w}, y_{w}) = 1 - \sum_{k=0}^{N} \binom{m}{k} \int_{0}^{k} \left( \frac{k_{2} x_{1}}{k_{1} \lambda_{M}} \right)^{p} e^{-\frac{x_{1}}{\lambda_{w}}} \left( \frac{k_{2} y_{w}}{\lambda_{M}} \right)^{k-p} dx \]
\[ (51) \]

Noticing that the RV \( y_{h_{w}} \) is decoupled with \( y_{w} \), the expectation can be taken separately. Hence, with the help of [22, Eq. (3.471.9)], we obtain
\[ \int_{0}^{\infty} e^{-\frac{k_{2} x}{\lambda_{w}}} \left( \frac{k_{2}}{k_{1} \lambda_{M}} \right)^{p} e^{-\frac{mN-1}{\lambda P}} \left( \frac{m}{\lambda P} \right)^{mN} \int_{0}^{k_{2}} e^{-\frac{mN-1}{\lambda P}} \left( \frac{m}{\lambda P} \right)^{mN} dx \]
\[ (52) \]

Similarly, invoking [22, Eq. (3.326.2)], we have
\[ \int_{0}^{\infty} e^{-\frac{k_{2} x}{\lambda_{w}}} \left( \frac{k_{2}}{k_{1} \lambda_{M}} \right)^{p} e^{-\frac{mN-1}{\lambda P}} \left( \frac{m}{\lambda P} \right)^{mN} dx = \frac{k_{2} x^{k-p}}{k_{1} \lambda_{M} \lambda_{P}} \] \[ (53) \]

To this end, pulling everything together yields the desired result.

\section*{APPENDIX B

PROOF OF PROPOSITION 1}

Starting from (49), conditioned on \( y_{w} \) and \( y_{w} \), we have
\[ P_{out}^{MRT}(R_{S}|y_{w}, y_{w}) = 1 - \left( 1 - \frac{y_{w}}{y_{m}} \right) \int_{k_{1}}^{\infty} \frac{x^{N-1}}{\left( \sum_{k=0}^{N} \binom{m}{k} x_{1}^{x_{1}} x_{2}^{x_{2}} \right)} e^{-\frac{x_{1} + x_{2}}{\lambda_{w}}} dx \]
\[ (54) \]

With the help of [22, Eq. (3.351.2)] and \( e^{-\frac{k_{2}}{\lambda_{w}}} = \sum_{k=0}^{\infty} \left( \frac{m}{\lambda_{w}} \right)^{k} \), conditioned on \( y_{w} \), the outage probability can be expressed as
\[ P_{out}^{MRT}(R_{S}|y_{w}) = 1 - \left( 1 - \frac{y_{w}}{y_{m}} \right) \sum_{k=0}^{\infty} \frac{1}{k!} \left( \frac{m}{\lambda_{w}} \right)^{k} \]
\[ (55) \]

Then, averaging over \( y_{w} \), with the help of [22, Eq. (3.471.9)], the secrecy outage probability can be computed as (56) on the top of next page

Expanding the Bessel function by [22, Eq. (8.446)] and omitting the high order items yield
\[ P_{out}^{MRT}(R_{S}) = \frac{1}{k!} \left( \frac{m}{k_{1} \lambda_{M} \lambda_{P}} \right)^{k} \Gamma(mN - k) \]
\[ (57) \]

By omitting the high order items, the desired result can be obtained.

\section*{APPENDIX C

PROOF OF THEOREM 2}

Based on (11) and (12), we set \( y_{m} = ||h_{m}||^{2} \) and \( y_{w} = ||h_{w}||^{2} \). In this case, since the selected antenna corresponds to a random transmit antenna for the eavesdropper, the pdf of \( y_{w} \) can be expressed as \( f_{y_{w}}(x) = \frac{1}{\lambda_{w}} e^{-\frac{x}{\lambda_{w}}} \). Now, with some simple algebraic manipulations, the cumulative distribution function (cdf) of \( y_{m} \) can be shown as
\[ F_{y_{m}}(x) = \left( \int_{0}^{x} e^{-\frac{t}{\lambda_{w}}} dt \right)^{N} = \left( 1 - e^{-\frac{x}{\lambda_{w}}} \right)^{N} \] \[ (58) \]
Hence the pdf of $y_{B,M}$ can be obtained with a simple derivative as follows:

$$f_{y_{B,M}}(x) = \frac{N}{\lambda M} \left(1 - e^{-\frac{x}{\lambda M}}\right)^{N-1} e^{-\frac{x}{\lambda M}}.$$  (59)

Then, the desired result can be obtained by following similar lines as in the proof of Theorem 1.

**APPENDIX D**

**PROOF OF THEOREM 3**

In this case, the strongest antenna corresponds to a random transmit antenna for the legitimate user, as such, the pdf of $y_{B,M}$ can be expressed as $f_{y_{B,M}}(x) = \frac{1}{\lambda M} e^{-\frac{x}{\lambda M}}$. Then, with some simple algebraic manipulations, the cdf of $y_{W}$ can be shown as

$$F_{y_{W}}(x) = 1 - \left(\int_{x}^{\infty} \frac{1}{\lambda W} e^{-\frac{t}{\lambda W}} dt\right)^{N} = 1 - e^{-\frac{x}{\lambda W}}.$$  (60)

Taking derivation of (60), the pdf of $y_{W}$ can be expressed as

$$f_{y_{W}}(x) = \frac{N}{\lambda W} e^{-\frac{x}{\lambda W}}.$$  (61)

Then, the desired result can be obtained by following similar lines as in the proof of Theorem 1.

**APPENDIX E**

**PROOF OF PROPOSITION 3**

By following the same steps as in the proof of Proposition 1, we obtain

$$P_{\text{out}}^{\text{TAS2}}(R_S) = 1 - \frac{N \lambda M}{N \lambda M + k_2 \lambda W} \times \frac{2}{\Gamma(mN)} \left(\frac{m(k_2 - 1)}{k_1 \lambda M \lambda P}\right)^{\frac{m}{2}} K_{m}(2) \left(\frac{m(k_2 - 1)}{k_1 \lambda M \lambda P}\right).$$  (62)

Expanding the Bessel function and omitting the high order items yield

$$P_{\text{out}}^{\text{TAS2}}(R_S) = 1 - \frac{N \lambda M}{N \lambda M + k_2 \lambda W} \times \left(1 - \frac{m(k_2 - 1)}{(mN - 1)k_1 \lambda P} \frac{1}{\lambda_M} + o\left(\frac{1}{\lambda_M}\right)\right).$$  (63)

Then, the desired result can be obtained along with some simple algebraic manipulations.

**APPENDIX F**

**PROOF OF THEOREM 4**

The exact secrecy outage probability is difficult to characterize, instead, we apply the following approximation as in [39]

$$C_S = \left[\log\left(\frac{1 + \gamma_M}{1 + \gamma_W}\right)\right]^{+} \approx \left[\log\left(\frac{\gamma_M}{\gamma_W}\right)\right]^{+} = \left[\log\left(\frac{|h_{M,k}|^2}{|h_{W,k}|^2}\right)\right]^{+}.$$  (64)

It is worth pointing out that, such approximation is reasonably tight and becomes asymptotically exact in the high SNR regime.

Then, we can define a new variable as follows

$$X = \max_{i=1,\cdots,N} \left(\frac{|h_{i,d}|^2}{|h_{i,c}|^2}\right).$$  (65)

Along with some simple algebraic manipulations, the cdf of $X$ can be computed as

$$F_{X}(x) = \left(\int_{0}^{x} \int_{\frac{1}{2}}^{\infty} \frac{1}{\lambda W} e^{-\frac{t}{\lambda W}} \frac{1}{\lambda M} e^{-\frac{t}{\lambda M}} dt dy\right)^{N} = \left(\frac{x}{x + \frac{1}{\lambda W}}\right)^{N}.$$  (66)

Taking derivation of (66), the pdf of $X$ can be expressed as

$$f_{X}(x) = N \frac{\lambda M}{\lambda W} \left(\frac{x}{x + \frac{1}{\lambda W}}\right)^{N-1} x^{N-1}.$$  (67)

To this end, the desired result can be expressed as $F_{X}(k_2)$.
as (69) on the top of the page.

The next step is to average over $y_{b|p}$. Here, we set

$$A = \frac{1}{\ln 2} \sum_{k=0}^{N-1} \left( \frac{1}{\lambda_M y_{b|p}} \right)^k \left( -\frac{1}{k!} \right)^k \frac{1}{\Gamma(mN) \lambda_{P}} \frac{m^{N+k}}{k!} \times
\int_{0}^{\infty} x^{mN-k-1} e^{-\frac{x}{k_{1M}x}} \Gamma \left( 0, \frac{1}{k_{1M}x} \right) dx.
$$

(70)

With the help of [22, Eq. (8.353.3)] and [22, Eq. (3.471.9)], (70) can be further expressed as

$$A = \frac{1}{\ln 2} \sum_{k=0}^{N-1} \left( \frac{1}{\lambda_M y_{b|p}} \right)^k \left( -\frac{1}{k!} \right)^k \frac{1}{\Gamma(mN) \lambda_{P}} \frac{m^{N+k}}{k!} \times
\int_{0}^{\infty} x^{mN-k-1} e^{-\frac{x}{k_{1M}x}} K_{mN-k} \left( 2x, \frac{m_{P}}{k_{1M}x} \right) dx.
$$

(71)

Invoking [22, Eq. (6.65.7)], we have

$$A = \sum_{k=0}^{N-1} \left( -\frac{1}{k!} \right)^k \frac{1}{\Gamma(mN) \lambda_{P}} \frac{m^{N+k}}{k!} \times
\frac{m_{N-k}}{k_{1M}x} S_{1-mN+k} \left( 2x, \frac{m_{P}}{k_{1M}x} \right).
$$

(72)

With the help of [22, Eq. (9.34.6)] and [22, Eq. (9.31.5)], we obtain

$$A = \frac{1}{\ln 2} \sum_{k=0}^{N-1} \left( -\frac{1}{k!} \right)^k \frac{1}{\Gamma(mN) \lambda_{P}} \frac{m^{N+k}}{k!} \times
G_{13}^{31} \left( \frac{m}{k_{1M}x_{P}}, 0 \right)
\left( 0, mN-k, 0, 0 \right).
$$

(73)

To this end, by following the same steps and utilizing [22, Eq. (3.326.2)], the desired result can be obtained.

**Appendix H**

**Proof of Proposition 6**

Capitalizing on the general framework proposed in [31] for the evaluation of asymptotic average secrecy rate, conditioned on $y_{b|p}$, we have

$$\bar{C}^{\infty} = A - B = \frac{1}{\ln 2} \int_{0}^{\infty} \ln(k_{1M}y_{b|p} x) f_{y_{b|p}}(x) dx - \frac{1}{\ln 2} \int_{0}^{\infty} \frac{k_{1M}y_{b|p}}{1 + k_{1M}y_{b|p} x} \left( 1 - F_{y_{b|w}}(x) \right) dx.
$$

(74)

With the help of [22, Eq. (4.352.1)], $A$ can be computed as

$$A = \log_{2}(\lambda_{M}) + \frac{1}{\ln 2} \left( \ln \frac{k_{1M}y_{p}}{m} + \psi(N) + \psi(mN) \right).$$

(75)

The calculation of $B$ is exactly the same as $C$ above. Then we have

$$B = \frac{1}{\ln 2} \int_{0}^{\infty} \ln \left( \frac{m_{N}}{k_{1M}x_{P}} \right) \frac{0}{mN, 0, 0} dx.
$$

(76)

The desired result then follows immediately.

**Appendix I**

**Proof of Theorem 8**

Starting from the definition, we have

$$C^{T_{A3}} = \int_{0}^{\infty} \ln(x) f_{X}(x) dx.
$$

(77)

Substituting (67) into (77) yields

$$C^{T_{A3}} = \frac{N_{W} \lambda_{W} m}{\ln 2} \int_{0}^{\infty} \ln x \frac{x^{N-1}}{(x + \frac{\lambda_{M}}{\lambda_{W}})^{N+1}} dx.
$$

(78)

Making a change of variable $t = x - 1$ and utilizing [40, Eq. (8.4.6.5)], (78) can be alternatively written as

$$C^{T_{A3}} = \frac{N \lambda_{W} m}{\ln 2} \lambda_{W} \int_{0}^{\infty} \ln(t+1) \frac{(t+1)^{N-1}}{(t+1 + \frac{\lambda_{M}}{\lambda_{W}})^{N+1}} dt.
$$

(79)

Applying the binomial expansion and utilizing [22, Eq. (7.811.5)] yield the desired result.

**Appendix J**

**Proof of Proposition 9**

By taking the general form, we obtain

$$\bar{C}^{\infty} = A - B = \frac{1}{\ln 2} \int_{0}^{\infty} \ln(x) f_{y_{b|w}}(x) dx - \frac{1}{\ln 2} \int_{0}^{\infty} \ln(y) f_{y_{b|w}}(y) dy,
$$

(80)

where the two pdfs $f_{y_{b|w}}(x)$ and $f_{y_{b|w}}(x)$ have been derived in [41] as

$$f_{y_{b|w}}(x) = \frac{N_{W} x^{N-1}}{\lambda_{M}^N} \Gamma(2 - N, \frac{x}{\lambda_{M}}),$$

(81)
and
\[
f_{SW}(x) = N(N-1) \sum_{k=0}^{N-2} \binom{N-2}{k} \frac{x^{k+1}}{k!} \Gamma \left( -1 - k, -\frac{x}{\lambda W} \right).
\]

For \( A \) utilizing [22, Eq. (8.350.2)] and [22, Eq. (4.352.1)], we obtain
\[
A = \frac{1}{\ln 2} \left( \ln \lambda M + \psi(2) - \frac{1}{N} \right).
\]

Following the similar lines, \( B \) can be computed as
\[
B = \frac{1}{\ln 2} \left( \ln \lambda W + \psi(1) + 1 - N + \sum_{k=2}^{N} (-1)^k \frac{N!}{k!} \right).
\]

The desired result then follows immediately.

REFERENCES


Xin Jiang (S’16) received the B.S. degree in information and communication engineering from Zhejiang University, Hangzhou, China, in 2015. He is currently working towards his M.S. degree in the college of information science and electronic engineering at Zhejiang University. His research interests include physical layer security and wireless powered communications.

Caijun Zhong (S’07-M’10-SM’14) received the B.S. degree in Information Engineering from the Xi’an Jiaotong University, Xi’an, China, in 2004, and the M.S. degree in Information Security in 2006, Ph.D. degree in Telecommunications in 2010, both from University College London, London, United Kingdom. From September 2009 to September 2011, he was a research fellow at the Institute for Electronics, Communications and Information Technologies (ECT), Queens University Belfast, Belfast, UK. Since September 2011, he has been with Zhejiang University, Hangzhou, China, where he is currently an associate professor. His research interests include massive MIMO systems, full-duplex communications, wireless power transfer and physical layer security.

Dr. Zhong is an Editor of the IEEE TRANSACTIONS ON WIRELESS COMMUNICATIONS, IEEE COMMUNICATIONS LETTERS, EURASIP JOURNAL ON WIRELESS COMMUNICATIONS AND NETWORKING, and JOURNAL OF COMMUNICATIONS AND NETWORKS. He is the recipient of the 2013 IEEE ComSoc Asia-Pacific Outstanding Young Researcher Award. He and his coauthors have been awarded a Best Paper Award at the WCSP 2013. He was an Exemplary Reviewer for IEEE TRANSACTIONS ON COMMUNICATIONS in 2014.

Xiaoming Chen (M’10-SM’14) received the B.Sc. degree from Hohai University in 2005, the M.Sc. degree from Nanjing University of Science and Technology in 2007 and the Ph.D. degree from Zhejiang University in 2011, all in electronic engineering. He is currently with the College of Electronic and Information Engineering, Nanjing University of Aeronautics and Astronautics, Nanjing, China. Since February 2015, he is a Humboldt Research Fellow at the Institute for Digital Communications, University of Erlangen-Nürnberg, Erlangen, Germany. His research interests mainly focus on multiple-antenna techniques, wireless security, interference network and wireless power transfer, etc.

Dr. Chen serves as an Associate Editor for the IEEE ACCESS and an Editor for the IEEE COMMUNICATIONS LETTERS. He was honoured as an Exemplary Reviewer of the IEEE Communications Letters in 2014 and the IEEE Transactions on Communications in 2015.

Zhaoyang Zhang (M’02) received his Ph.D. degree in communication and information systems from Zhejiang University, China, in 1998. He is currently a full professor with the Department of Information Science and Electronic Engineering, Zhejiang University. His research interests are mainly focused on information theory and coding theory, signal processing for communications and in networks, and their applications in the next generation wireless mobile communication systems. He has co-authored more than 150 refereed international journal and conference papers as well as two books in the above areas. He was a co-recipient of several conference Best Paper Awards / Best Student Paper Award. He is currently serving as Editor for IEEE TRANSACTIONS ON COMMUNICATIONS, IET COMMUNICATIONS and some other international journals. He has served or is serving as TPC Co-Chair or Symposium Co-Chair for many international conferences like WCSP 2013, ICUFN2011/12/13, and Globecom 2014 Wireless Communications Symposium, etc.