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Impact of Correlated Fading Channels on Multiple Secure Relaying with Outdated Relay Selection

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Abstract—In this paper, we study the impact of correlated fading channels on multiple secure decode-and-forward (DF) relaying with outdated relay selection, where the information transmission assisted by the \( N \) DF relays from the source to the destination can be overhead by the eavesdropper in the network. The eavesdropping channels are correlated with the main channels, which affects the network security. To enhance the network security, one best relay is chosen to assist the secure transmission, which is however maybe outdated in time-varying channel environments. The impact of both channel correlation and outdated relay selection on the secrecy performance is studied by deriving the analytical expression of the secrecy outage probability (SOP). The asymptotic SOP is also provided with high main-to-eavesdropper ratio (MER). From the asymptotic SOP, we find that only the outdated degree of relay selection affects the network secrecy diversity order, but the channel correlation does not. Moreover, it is interesting to find that the channel correlation is beneficial to the transmission security in the high MER regime.

Index Terms—Secure communications, correlated fading channels, outdated relay selection, secrecy diversity order.

I. INTRODUCTION

Due to the broadcast nature, the wireless transmission may be overhead by eavesdroppers in the network, and the severe issue of information leakage arises [1]. To prevent the wiretap, the encryption algorithm and physical-layer security (PLS) have been studied in the literature. The pioneering work of PLS is investigated by Wyner [2], where the wiretap channel model was firstly proposed. Then researchers extended to fading channels and studies the important metrics of secrecy performance, such as secrecy capacity and secrecy outage probability (SOP) [3]–[5]. As relaying is a promising technique for the next-generation communications, it is of vital importance to investigate the PLS of relay networks. For amplify-and-forward (AF) and decode-and-forward (DF) relaying, the secrecy performance has been studied by deriving the analytical expression of SOP [6]–[8]. Moreover, the asymptotic SOP with high main-to-eavesdropper ratio (MER) was provided to obtain the insights on the system.

In the most existing works of the literature, the eavesdropping channels are assumed to be independent of the main channels. However, this ideal assumption may not hold in practice due to many factors such as antenna deployments, proximity of the legitimate receiver and eavesdropper, and scattering environments. The impact of channel correlation between the main and eavesdropping links on the secrecy performance is studied in [9], [10], where it has been found that the channel correlation is harmful to the transmission security in the low MER region. To enhance the transmission security, the opportunistic selection is an effective technique to exploit the channel fluctuation between antennas, users and relays [7], [11]. The selection can be implemented in a centralized or distributed manner through dedicated feedback links, which may take some time to complete. In time-varying channel environments, the channels may vary from the instant of selection to that of actual data transmission, which causes the selection based on the outdated channel state information (CSI) [12] and severely limits the system secrecy performance [13], [14].

In this paper, we investigate the secure multi-relay networks with channel correlation between the main and eavesdropping links, where the information transmission assisted by \( N \) DF relays from the source to the destination may be overhead by the eavesdropper in the network. To strengthen the secure transmission, one best relay is selected among \( N \) relays, which is however based on the outdated CSI in time-varying channel environments. We study the impact of both channel correlation and outdated relay selection by deriving the analytical and asymptotic expressions for the SOP. From the asymptotic result, we find that the secrecy diversity order is equal to \( N \) only in perfect channel state information; otherwise it degenerates to unity. Moreover, the channel correlation does not affect the secrecy diversity order, but it can help strength the secure transmission in high MER region.

Notations: The notation \( \mathcal{C} \mathcal{N}(0, \sigma^2) \) denotes a circularly symmetric complex Gaussian random variable (RV) with zero mean and variance \( \sigma^2 \). We use \( f_X(\cdot) \) to represent the probability density function (PDF) of the RV \( X \). In addition, \( I_0(x) \) is the modified Bessel function of the first kind of order zero [15], \( P_2(\cdot) \) returns the probability, and \( E[\cdot] \) denotes the statistical average. We use \( h_{A,B} \) to denote the channel parameter of the \( A-B \) link.

II. SYSTEM MODEL

Fig. 1 shows the system model of a two-phase secure multiple DF relays network, where the information transmission...
from the source S to the destination D can be overheard by the eavesdropper E. There are no direct links from S to D and E, and the information transmission are only through the N relays \( \{ R_n \}_{1 \leq n \leq N} \). One best relay is chosen among N available relays to assist the secure transmission. However, the selection may be outdated in time-varying channel environments, and the best relay is not always chosen. In addition, due to the reasons such as antenna deployment and radio scattering, the channels at the receivers D and E are correlated with each other, i.e., \( h_{R_n,D} \) is correlated with \( h_{R_n,E} \). Due to the size limitation, the nodes in the network are equipped with a single antenna, and all links experience time-varying Rayleigh fading. In the following, we present the two-phase secure data transmission process and the relay selection criterion for the considered system.

Suppose that the relay \( R_n \) is used for the two-phase secure data transmission. In the first phase, the source S sends signal \( x_S \) to \( R_n \) with transmit power \( P \), and \( R_n \) receives

\[
y_R = \sqrt{P}h_{S,R_n}x_S + n_{R_n},
\]

where \( h_{S,R_n} \sim CN(0, \alpha) \) denotes the channel parameter of the S–\( R_n \) link, and \( n_{R_n} \sim CN(0, \sigma^2) \) is the additive white Gaussian noise (AWGN) at the relay \( R_n \). If the relay \( R_n \) can correctly decode the message from the source, i.e., it can support a target data rate \( R_t \),

\[
\frac{1}{2} \log_2 \left( 1 + \frac{P|h_{R_n,S}|^2}{\sigma^2} \right) \geq R_t,
\]

the relay forwards the message to D in the second phase with the transmit power \( P \). Accordingly, D and E receive

\[
\begin{align*}
y_D &= \sqrt{P}h_{R_n,D}x_S + n_D, \\
y_E &= \sqrt{P}h_{R_n,E}x_S + n_E,
\end{align*}
\]

where \( h_{R_n,D} \sim CN(0, \beta) \) and \( h_{R_n,E} \sim CN(0, \varepsilon) \) are the channel parameters of the \( R_n \)–D and \( R_n \)–E links, respectively. The noise terms \( n_D \sim CN(0, \sigma^2) \) and \( n_E \sim CN(0, \sigma^2) \) are the AWGN at D and E, respectively.

We use \( u = |h_{S,R_n}|^2 \), \( v = |h_{R_n,D}|^2 \) and \( w = |h_{R_n,E}|^2 \) to represent the channel gains of the S–\( R_n \), \( R_n \)–D and \( R_n \)–E links, respectively. The secrecy outage occurs when the data rate difference between the main and eavesdropping links falls below a target secrecy data rate \( R_s \),

\[
\frac{1}{2} \log_2 \left( 1 + \frac{P}{\sigma^2} v \right) - \frac{1}{2} \log_2 \left( 1 + \frac{P}{\sigma^2} w \right) < R_s,
\]

which is equivalent to

\[
\frac{1 + P v_n}{1 + P w_n} < \gamma_s,
\]

where \( \gamma_s = 2^{2R_s} \) is the secrecy SNR threshold.

To enhance the transmission security, we need to select one best relay to strength the secure data transmission. Let \( \Omega \) denote the candidate set of relays that can successfully decode the message from the source. Then based on the main channels only\(^1\), the relay selection is performed to choose one best relay \( R_n^* \) among \( \Omega \),

\[
n^* = \arg \max_{n \in \Omega} v_n,
\]

which maximizes the received SNR at the destination D.

### III. Channel Correlation and Outdated Relay Selection

In this work, we consider the correlated channels between the receivers D and E, and the correlation between \( v_n \) and \( w_n \) is characterized by [10]

\[
f_{w_n|v_n}(w|v) = I_0 \left( \frac{2}{1-\rho_c} \sqrt{\frac{\rho_{vw} \sqrt{\rho_c}}{\beta \varepsilon}} \right) e^{-\frac{\sqrt{\rho_{vw}} v + w}{1-\rho_c}},
\]

where \( \rho_c \in [0, 1] \) is the power correlation coefficient. Specifically, \( \rho_c = 0 \) represents that \( v_n \) is independent of \( w_n \), while \( \rho_c = 1 \) denotes the completely linear correlation.

Besides the channel correlation, the outdated relay selection has a significant impact on the network security. The selection of (7) can be implemented in a distributed or centralized manner in practice, through some dedicated feedback channels [7]. However, due to the limited feedback resources, the channels may vary from the instant of relay selection to that of actual data transmission in time-varying channel environments. Let \( v_{n^*} \) and \( v_{n^*} \) denote the channels of \( R_{n^*} \)–D at the instants of relay selection and actual data transmission, respectively. The outdated selection can be characterized by the conditional PDF \( f_{v_{n^*}|v_n}(v|\tilde{v}) \) [12],

\[
f_{v_{n^*}|v_n}(v|\tilde{v}) = \frac{1}{(1-\rho_o)\beta} e^{-\frac{\rho_{ov} v + \tilde{v}}{(1-\rho_o)\beta}} I_0 \left( \frac{2\sqrt{\rho_{ov} v}}{(1-\rho_o)\beta} \right),
\]

where \( \rho_o \in [0, 1] \) denotes the outdated degree of relay selection. In particular, \( \rho_o = 1 \) denotes that the selection is based on the perfect CSI, while \( \rho_o = 0 \) represents the completely outdated relay selection.

In the following, we will study the impact of channel correlation and outdated relay selection on the network secrecy performance by providing the analytical and asymptotic expressions of secrecy outage probability.

\(^1\)The instantaneous channel parameters of eavesdropping links are not involved in the relay selection criterion, since they are generally hard to obtain in practice.
\[ P_{\text{out}} = \sum_{K=1}^{N} \binom{N}{K} e^{-\frac{Kn}{P_v}} (1 - e^{-\frac{kn}{P_v}})^{N-K} \left( 1 - \sum_{k=1}^{K} \sum_{m=0}^{T} \sum_{i=0}^{m} \sum_{j=0}^{m} C_{km} d_{m,ij} (m+j)! (b_1 \gamma_s + b_2)^{-(m+j+1)} \right). \]  
\[ (25) \]

IV. SECRECY OUTAGE PROBABILITY

A. Analytical Expression

Note that the set \( \Omega \) may have \( K(K = 1, 2, \ldots, N) \) candidates that can correctly decode the message from the source, we can express the SOP of the considered system as

\[ P_{\text{out}} = \sum_{K=1}^{N} \binom{N}{K} \Pr \left[ \tilde{P}_{u1} \geq \gamma_t, \cdots, \tilde{P}_{uK} \geq \gamma_t, \right. \]
\[ \left. \tilde{P}_{uK+1} < \gamma_t, \cdots, \tilde{P}_{uN} < \gamma_t \right\rfloor \frac{1 + \tilde{P}_{v_{n^*}}}{1 + \tilde{P}_{w_{n^*}}} < \gamma_s \right], \]  
\[ (10) \]

where \( \gamma_t = 2^{2R_t} - 1 \) denotes the SNR threshold of successful decoding at the relay. Since the random variable \( u_{n^*} \) is independent of \( v_{n^*} \) and \( w_{n^*} \), we can rewrite \( P_{\text{out}} \) as

\[ P_{\text{out}} = \sum_{K=1}^{N} \binom{N}{K} \Pr \left[ \tilde{P}_{u1} \geq \gamma_t, \cdots, \tilde{P}_{uK} \geq \gamma_t, \right. \]
\[ \left. \tilde{P}_{uK+1} < \gamma_t, \cdots, \tilde{P}_{uN} < \gamma_t \right\rfloor \frac{1 + \tilde{P}_{v_{n^*}}}{1 + \tilde{P}_{w_{n^*}}} < \gamma_s \right], \]  
\[ (11) \]

where the PDF of \( f_{u_n}(u) = \frac{1}{\alpha} e^{-\frac{u}{\alpha}} \) is applied and the probability \( J_K \) denotes the conditional SOP with the given \( K \) active relays. From the selection criterion in (7) and the conditional PDF of \( f_{v_{n^*}}(v|v) \) in (9), the PDF of \( v_{n^*} \) is given by [12]

\[ f_{v_{n^*}}(v) = \sum_{k=1}^{K} \binom{K}{k} \frac{k(-1)^{k-1}}{\beta^k [k(1 - \rho_o) + \rho_o]} e^{-\frac{k(v - \rho_c \gamma_s)}{1 - \rho_c \gamma_s}}. \]  
\[ (13) \]

From the above \( f_{v_{n^*}}(v) \) and the conditional \( f_{w_n|v_n}(w|v) \) in (8), we obtain the joint PDF of \( v_{n^*} \) and \( w_{n^*} \) as,

\[ f_{v_{n^*}, w_{n^*}}(v, w) = f_{w_n|v_n}(w|v) f_{v_{n^*}}(v) \]
\[ = \sum_{k=1}^{K} \binom{K}{k} \frac{k(-1)^{k-1}}{\beta \varepsilon [k(1 - \rho_o) + \rho_o]} f_0 \left( 1 - \rho_c \varepsilon e^{-\frac{k w}{\beta \varepsilon}} \right) e^{-\frac{\rho_c \gamma_s}{1 - \rho_c \gamma_s}} \gamma_s^{\frac{\rho_c \gamma_s}{1 - \rho_c \gamma_s} - 1 - j}, \]
\[ (14) \]

Note that the Bessel function \( I_0(x) \) can be expanded by the series as [15]

\[ I_0(x) = \sum_{m=0}^{\infty} \frac{x^{2m}}{4^m (m!)^2} \]
\[ \approx \sum_{m=0}^{T} \frac{x^{2m}}{4^m (m!)^2}, \]  
\[ (16) \]

where the truncation error in (16) decays exponentially with \( T \) [10]. Hence with an efficient number of series, we can obtain an accurate approximation for \( I_0(x) \). In the following, we ignore the approximation error, and re-express \( f_{v_{n^*}, w_{n^*}}(v, w) \) as

\[ f_{v_{n^*}, w_{n^*}}(v, w) = \sum_{k=1}^{T} \sum_{m=0}^{T} C_{km} v^{m} w^{m} e^{-b_1 v} e^{-b_2 w}, \]  
\[ (17) \]

with

\[ C_{km} = \binom{K}{k} \frac{k(-1)^{k-1}}{\beta \varepsilon [k(1 - \rho_o) + \rho_o]} \frac{\rho_c^m}{(1 - \rho_c \gamma_s)^{m+1}}, \]  
\[ (18) \]

\[ b_1 = \frac{k}{k(1 - \rho_o) + \rho_o + 1 - \rho_c \gamma_s}, \]  
\[ (19) \]

\[ b_2 = \frac{1}{(1 - \rho_c \gamma_s)^2}. \]  
\[ (20) \]

From (17), we can compute \( J_K \) as

\[ J_K = \Pr \left( v_{n^*} < \gamma_s - \frac{1}{P} + \gamma_s w_{n^*} \right) \]
\[ = \int_0^{\infty} \int_0^{\gamma_s w} f_{v_{n^*}, w_{n^*}}(v, w) dv dw \]
\[ = 1 - \sum_{k=1}^{K} \sum_{m=0}^{T} \sum_{i=0}^{m} \sum_{j=0}^{m} C_{km} d_{m,ij} (m+j)! \times (b_1 \gamma_s + b_2)^{-(m+j+1)}, \]  
\[ (23) \]

with

\[ d_{m,ij} = \frac{m!}{i!} \left( \frac{i}{j} \right) e^{-b_1 (\gamma_s - 1)/P} \frac{\gamma_s^j (\gamma_s - 1)^{-j}}{P^i}. \]  
\[ (24) \]

By applying the result of \( J_K \) into (12), we can obtain the analytical SOP for the multiple secure relaying with channel correlation and outdated relay selection, as shown in (25) at the top of this page. Note that the obtained analytical SOP consists of elementary functions only, and hence is easily to be evaluated.

B. Asymptotic Expression

To obtain the insights on the system, we now extend to derive the asymptotic SOP with high MER. With a large transmit power \( P \), we can approximate \( P_{\text{out}} \) as

\[ P_{\text{out}} \approx J_N \]
\[ \approx \Pr \left( \frac{v_{n^*}}{w_{n^*}} < \gamma_s \right). \]  
\[ (27) \]
Let \( z = \frac{2v_n - n^*}{w_n n^*} \), and then the PDF of \( z \) is derived as

\[
f_z(z) = \int_0^\infty w f_{v_n^*, w_n^*}(wz, w) dw
\]

\[
= \sum_{k=1}^N \binom{N}{k} \frac{(-1)^{k-1}k}{[k(1-\rho_c) + \rho_o]^{k-1}} \left( \frac{z}{[k(1-\rho_c) + \rho_o]} + \frac{\rho_o}{\rho_c} + \frac{1}{\rho_c} \right) \frac{1}{[k(1-\rho_c) + \rho_o]^{k-1}},
\]

where \( \lambda = \beta / \varepsilon \) is the MER. From the above \( f_z(z) \), we then approximate \( P_{\text{out}} \) as

\[
P_{\text{out}} \approx \int_{-\infty}^{\gamma_0} f_z(z) dz
\]

\[
\approx \frac{N!}{\lambda} \left( \frac{1-\rho_c}{\gamma_0} \right)^N, \quad \text{If } \rho_o = 1
\]

\[
\approx \left( \frac{1-\rho_c}{\gamma_0} \right) \left( \sum_{k=1}^N \binom{N}{k} \frac{k(1-\rho_c)}{k(1-\rho_c) + \rho_o} \right), \quad \text{If } \rho_o < 1
\]

Fig. 2 demonstrates the numerical and simulated secrecy outage probabilities versus the transmit power \( P \), where \( N = 3, \lambda = 20 \text{dB} \) and \( D = 0.5 \). Several cases of channel correlation are considered with \( \rho_c = 0.3, 0.5 \) and 0.7. Both perfect and outdated CSI environments are studied with \( \rho_o = 1.0 \) and 0.5, respectively. As observed from the figure, for different values of \( \rho_c \) and \( \rho_o \), the analytical result matches well with the simulation one, and the asymptotic result converges to the exact one when \( P \) is large. This validates the derived analytical and asymptotic expressions of the SOP. Moreover, in both perfect and outdated CSI environments, the secrecy performance improves with larger \( \rho_o \), as better CSI helps select the best relay to assist the secure transmission. Furthermore, the secrecy performance improves with larger \( P \). The improvement is saturated in high \( P \) region, as the fixed MER becomes the bottleneck of the secrecy performance.

Fig. 3 illustrates the secrecy outage probability versus MER with different values of \( \rho_c \) and \( \rho_o \), where \( P = 40 \text{dB} \) and \( N \) varies from 1 to 3. We can see from this figure that for different values of \( N, \rho_c \) and \( \rho_o \), the analytical result fits well with the simulation one, and the asymptotic result converges to the simulation one in the high MER region. Moreover,

Note that in the low MER region, the channel correlation is however harmful to the secure transmission [9], [10].

V. NUMERICAL AND SIMULATION RESULTS

In this section, we provide some numerical and simulation results to verify the proposed studies for the multiple secure relaying with channel correlation and outdated relay selection. All links in the network experience Rayleigh fading, and we adopt the path-loss model with exponent of four to measure the average channel gains of main links. Without loss of generality, we normalize the distance between the source and destination to unity, where the relays are in between of them. Let \( D \) denote the distance between the source and relays, and accordingly, \( \alpha = D^{-4} \) and \( \beta = (1 - D)^{-4} \) are set. The target data rate \( R_t \) of the first hop is 1 bps/Hz, and hence the associated \( \gamma_t \) is 3. The secrecy data rate \( R_s \) is 0.2 bps/Hz, so that the secrecy SNR threshold \( \gamma_s \) is 1.32.

As observed from the figure, for different values of \( \rho_c \) and \( \rho_o \), the analytical result matches well with the simulation one, and the asymptotic result converges to the exact one when \( P \) is large. This validates the derived analytical and asymptotic expressions of the SOP. Moreover, in both perfect and outdated CSI environments, the secrecy performance improves with larger \( \rho_o \), as better CSI helps select the best relay to assist the secure transmission. Furthermore, the secrecy performance improves with larger \( P \). But the improvement is saturated in high \( P \) region, as the fixed MER becomes the bottleneck of the secrecy performance.
only the value of \( \rho_o \) affects the curve slopes, but \( \rho_c \) does not. This indicates that the secrecy diversity order depends on the outdated degree of relay selection, but not on the channel correlation. In particular, the curves are in parallel with each other for different relay numbers when \( \rho_o = 0.5 \), indicating that the system secrecy diversity order is unity for different values of \( \rho_c \), as long as the selection is outdated. On the contrast, the curve slope is proportional to \( N \) in the perfect CSI with \( \rho_o = 1 \), indicating that the system full secrecy diversity order can be achieved for different values of \( \rho_c \). Such observations validate the insights from the asymptotic expression of the SOP.

Fig. 3 depicts the impact of \( \rho_c \) and \( \rho_o \) on the secrecy outage probability with \( N = 3 \) and \( \lambda = 20 \text{dB} \), where \( \rho_o \) ranges from 0 to 1 and \( \rho_c \) varies in \{0, 0.3, 0.5, 0.7, 0.9\}. We can find that for different values of \( \rho_c \) and \( \rho_o \), the analytical result matches well with the simulation one. Moreover, the secrecy performance improves with larger \( \rho_c \) and \( \rho_o \). This is because that higher channel correlation helps the destination have more information about the channel fluctuation of eavesdropping links, and better CSI helps select the best relay to assist the secure transmission.

VI. Conclusion

In this paper, we investigated the secure relaying with \( N \) DF relays, where the main and eavesdropping channels are correlated. The relay selection was performed to choose one best relay to assist the secure transmission, which is however maybe outdated. We studied the impact of channel correlation and outdated relay selection on the secrecy performance by deriving the analytical SOP as well as the asymptotic expression with high MER. From the asymptotic SOP, we found that the channel correlation does not affect the secrecy diversity order, but can be beneficial to the transmission security. Importantly, we confirmed that the full secrecy diversity order of \( N \) is achieved only when the selection is based on the perfect CSI.

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