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Published in:
IEEE Transactions on Signal Processing

Document Version:
Peer reviewed version

Queen's University Belfast - Research Portal:
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Download date:09. Nov. 2022
Multi-user regularized zero-forcing beamforming

L. D. Nguyen, H. D. Tuan, T. Q. Duong, and H. V. Poor

Abstract—Regularized zero-forcing beamforming (RZFB) is an interesting class of linear signal processing problems, which is very attractive for use in large-scale communication networks due to its simple visualization as a straightforward extension of the well-accepted zero-forcing beamforming (ZFB). However, unlike ZFB, which is multi-user interference free, RZFB must manage multi-user interference to achieve its high throughput performance. Most of existing works focus on the performance analysis of particular RZBF schemes such as the equip-power allocated RZBF under a fixed regularization parameter. This paper is the first work to consider the joint design of power allocation and regularization parameter for RZFB to maximize the worst users’ throughput or the quality-of-service awarded energy efficiency under a fixed transmit power constraint. Such designs pose very computationally challenging optimization problems, for which the paper proposes two-stage optimization algorithms of low computational complexity. Their computational and performance efficiencies are substantiated through numerical examples.

Index Terms—Multi-antenna communications, regularized zero-forcing beamforming, nonconvex optimization algorithms

I. INTRODUCTION

Linear beamforming is considered as the main driver of 5G signal processing techniques [1]. Zero-forcing beamforming (ZFB) [2] is a very popular class of linear beamforming due to its simplicity in design [3], [4]. Based on the channel matrix right inversion, ZFB is most attractive when the number of transmit antennas is much larger than the number of users (UEs) as the case of massive multiple input multiple output (MIMO) serving a few UEs [5]. Furthermore, ZFB for heterogenous networks has been considered in [6] and [7]. Reference [8] presented a multi-cell edge-aware ZFB to suppress interference at cell-edge UEs, where the main results are derivations for the number of UEs and antennas going to infinity. Reference [9] considered the two-stage beamformer design for massive MIMO to maximize the signal-to-leakage-plus-noise-ratio, where ZFB is used to eliminate the intra-group user interference. ZFB is often based on equal-power allocation to all UEs for computational tractability, which achieves the performance that is far from its capacity [10].

This work was supported in part by Institute for Computational Science and Technology, Hochiminh city, Vietnam, in part by the Australian Research Council’s Discovery Projects under Project DP190102501, in part by the U.K. Royal Academy of Engineering Research Fellowship under Grant RF1415\_14\_22, and in part by U.S. National Science Foundation under Grants CCF-093970 and CCF-1513915.

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When the number of transmit antennas is not much larger than the number of served UEs, the channel right inversion may not exist or is ill-posed, making ZFB either useless or poorly perform. A practical remedy for the channel matrix inversion is regularized zero-forcing beamforming (RZFB) [11]–[13], which uses a regularization parameter to make the channel matrix well-posed for the right inversion. Unlike ZFB, multi-user interference cannot be cancelled by RZFB and thus needs to be appropriately managed by optimizing the regularization parameter. So far, most of the existing works considered the problem of optimizing the regularization parameter based on the equal-power allocation to UEs, for which the regularization parameter is analytically found for the case of sum rate capacity under sum power constraint, providing that both number of transmit antennas and number of served UEs go to infinity with fixed ratio [11]. This value of the regularization parameter is used for analyzing the spectral efficiency of RZFB [14] or comparing the performance of RZFB with that of other schemes [15]. An important contribution is [16], which derived the deterministic equivalence of signal-to-interference-plus-noise (SINR) in the sense of almost sure probability when both the numbers of antennas and UEs go to infinity with a bounded ratio. The analytical formulas in [16, Corollaries 1 and 2] indeed rule out the opportunity of power allocation optimization as the individual rates are the same function of their allocated power. A large scale analysis has been made in [17], [18] under setting the regularization parameter to one. Ref [19] used a polynomial expansion for reducing the complexity of matrix inversion calculation, which works for SINR large scale analysis. An optimal ratio of the numbers of UEs and antennas for maximizing the asymptotic signal-to-leakage-plus-noise-ratio when both of these numbers go to infinity has been derived in [20]. Again, a heuristic method of low complexity for MIMO RZFB at high signal-to-noise ratio (SNR) was proposed in [21] with achievable sum rate analysis. There is no constraint for quality-of-service (QoS) for UEs in these works, making the multi-user services less meaningful.

Motivated by the above aforementioned existing issues in RZFB design, in this paper considers the joint design of power allocation and regularization parameter to maximize the worst UEs’ throughput or the energy-efficiency subject to QoS constraints for UEs in terms of their throughput. Under a fixed regularization parameter, such problems for single-antenna UEs have been particularly considered in our previous work [22]. However, the joint optimization in the power allocation and regularization parameter causes much more computational challenges that make the computational method [22] no longer applicable, even for the case of single-antenna UEs considered there. Our contribution is thus two-fold:

- Development of computational procedures for RZFB for
single-antenna UEs, which constitute of two stages. As the optimization problems are very complicated, in the first stage path-following computational procedures are developed to solve their nonconvex relaxed problems, which are simpler than the original problem but still not convex. In the second stage, exact penalty optimization algorithms, which compensate the relaxation done in the first stage, are employed to locate a stationary point of the original problems.

- Development of two-stage computational procedures for RZFB for multi-antenna UEs, which are very different from that for single-antenna UEs in term of mathematical structures and thus require quite different mathematical tools for their computation.

The paper is organized as follows. After this Introduction, Section II is devoted to RZFB optimization for single-antenna UEs while Section III is devoted to RZFB optimization for multi-antenna UEs. Numerical examples are used to demonstrate the viability of the Algorithms developed in the previous section in Section IV. Conclusions are drawn in Section V. The two appendices serve as the mathematic foundations for the developments in Section II and Section III. To the authors’ best knowledge, these mathematic results are new.

**Notation.** Bold-faced lower-case and upper-case letters, e.g., $\mathbf{x}$ and $\mathbf{X}$, are respectively used for vectors and matrices, while lower-case letters, e.g., $x$, are used for scalars. The inner product between the vectors $\mathbf{x}$ and $\mathbf{y}$ is defined as $\langle \mathbf{x}, \mathbf{y} \rangle = \mathbf{x}^H \mathbf{y}$. Analogously, $\langle \mathbf{X}, \mathbf{Y} \rangle = \text{trace}(\mathbf{X}^H \mathbf{Y})$ for the matrices $\mathbf{X}$ and $\mathbf{Y}$. The shorthand notation $\mathbf{X}^\dagger$ for the matrix $\mathbf{X}$ denotes the Hermitian symmetric positive definite matrix $\mathbf{X} \mathbf{X}^H$. $\mathbf{I}_N$ stands for the identity matrix of $N$ dimensions. The notation $\mathbf{A} \succ 0$ means that $\mathbf{A}$ is a Hermitian symmetric strictly positive definite. $\text{diag}[\alpha]_{i=1}^N \equiv \{\langle x_1, \ldots, x_N \rangle^T \in \mathbb{R}^N : x_i > 0, i = 1, \ldots, N\}$.

II. MISO REGULARIZED ZERO-FORCING BEAMFORMING

Consider a communication system with one base station (BS) serving $N$ UEs. The BS is equipped with a large-scale $M$-antenna array while the UEs are equipped with a single antenna. The multi-input single-output (MISO) equation of such system is

$$\mathbf{y} = \text{diag}[\sqrt{\beta_n}]_{n=1}^N \mathbf{H}^H \mathbf{s} + \mathbf{w},$$

where

$$\mathbf{H}^H = [h_{ij}^*]_{i=1,\ldots,N;j=1,\ldots,M} = \begin{bmatrix} h_{11}^H & \cdots & h_{1M}^H \\ \vdots & \ddots & \vdots \\ h_{N1}^H & \cdots & h_{NM}^H \end{bmatrix} \in \mathbb{C}^{N \times M},$$

and $\mathbf{y} \in \mathbb{C}^N$ is the signal received at $N$ UEs, $\mathbf{s} \in \mathbb{C}^M$ is the signal sent from BS and $\mathbf{w} \in \mathbb{C}^N(\sigma^2 \mathbf{I}_N)$ is the background noise. The channel from the BS to user $n$ is thus

$$\sqrt{\beta_n} \mathbf{h}_n, \mathbf{h}_n = [h_{n1} \ldots h_{nM}]^T,$$

where $\sqrt{\beta_n}$ expresses the large-scale fading, while $\mathbf{h}_n = \mathbf{R}_n^{1/2} [h_{n1} \ldots h_{nM}]^T$ with $h_{ij} \in \mathcal{CN}(0, 1)$ expressing the small scale fading and $\mathbf{R}_n$ expressing the spatial correlation due to netted spacing of large numbers of antennas. Like the existing works on beamforming such as [11]–[13], [23]–[26], we assume the availability of the full channel state information and refer the reader to [27] and references therein for the mechanism of large-scale MIMO channel estimation.

In the above system, unlike the conventional massive MIMO, we do not assume that $N \ll M$, under which the right inversion for $\hat{\mathbf{H}} = \text{diag}[\sqrt{\beta_n}]_{n=1}^N \mathbf{H}$ may not exist or be ill-conditioned; thus ZFB is not applicable. We are interested in the following class of RZFB, which admits a large-scale tractable computation

$$\hat{\mathbf{s}} = \mathbf{F} \mathbf{s} = \mathbf{H}(\mathbf{H}^H \mathbf{H} + \alpha \mathbf{I}_N)^{-1} \text{diag}[\sqrt{\beta_n}]_{n=1}^N \mathbf{s},$$

where $\mathbf{s} = (s_1, \ldots, s_N)^T$ is the information intended for the UEs, $\alpha > 0$ is the scalar to regularize $\mathbf{H}^H \mathbf{H}$, and $p_n, n = 1, \ldots, N$ are the power control coefficients. Instead of designing the complex beamforming matrix $\mathbf{F}$ in (3) of the large size $M \times N$, which obviously leads to a large scale intractable computation (see e.g. [23], [28]), by (3) we aim to design the real $N$-dimensional vector $\mathbf{p} \triangleq (p_1, \ldots, p_N)^T \in \mathbb{R}^N$ of power allocation for computational tractability. By (3) we propose a new RZFB compared to the existing RZFB defined by [11]–[13]

$$\mathbf{F} = \hat{\mathbf{H}}(\mathbf{H}^H \hat{\mathbf{H}} + \alpha \mathbf{I}_N)^{-1} \text{diag}[\sqrt{\beta_n}]_{n=1}^N.$$
while by using identity (7), the transmit power at the BS is expressed by

$$\pi(p, \alpha) = \text{trace}(F^H F) = \sum_{n=1}^{N} p_n (H_n^H D(\alpha))^2 H_n$$  \quad (12)$$
as shown by Appendix III.

We consider the following max-min throughput optimization problem

$$\max_{\mathbf{p} \in \mathbb{R}_+^N, \alpha > 0} \min_{n=1, \ldots, N} r_n(p, \alpha)$$  \quad (13a)

subject to

$$\pi(p, \alpha) \leq P, \quad (13b)$$

where $P$ is a given power threshold, i.e., we seek to maximize the worst UEs’ throughput under the total transmit power constraint.

We also consider the following quality-of-service (QoS)-aware energy efficiency problem [29], [30]

$$\max_{\mathbf{p} \in \mathbb{R}_+^N, \alpha > 0} \phi(p, \alpha) \triangleq \frac{\sum_{n=1}^{N} r_n(p, \alpha)}{\eta \pi(p, \alpha) + P_{\text{non}}}$$  \quad (14a)

subject to

$$r_n(p, \alpha) \geq \bar{r}, n = 1, \ldots, N, \quad (14b)$$

with a throughput threshold $\bar{r}$ to express the UEs’ QoS. Here $\eta > 1$ is reciprocal of the drain efficiency of the amplifier of BS. $P_{\text{non}} = M P_a + P_c$ is the circuit non-transmit power consumed at BS where $P_a$ and $P_c$ are circuit power per antenna and non-transmission power. The denominator of the objective function in (14) represents the total power consumption in broadcasting signals to UEs.

One can see from the definition (11) and (12) that the throughput function $r_n(p, \alpha)$ is a complex nonconcave function, while the sum power function $\pi(p, \alpha)$ is nonconvex. Therefore, both the max-min throughput optimization problem (13) and the energy-efficiency maximization problem (14) are maximization of a nonconcave objective function over nonconvex constraints, which are very computationally difficult. Our next subsections are devoted to their computation.

### A. Max-min throughput optimization

To address the max-min throughput optimization problem (13), make the following variable changes

$$x_n = 1/p_n > 0, n = 1, \ldots, N \quad (15)$$
to express the rate function $r_n$ and sum power function $\pi$ as the following functions of $x = (x_1, \ldots, x_N) \in \mathbb{R}_+^N$ and $\alpha > 0$:

$$r_n(p, \alpha) \rightarrow \tilde{r}_n(x, \alpha) \triangleq \ln \left( 1 + \frac{(H_n^H D(\alpha)H_n)^2}{x_n \left( \sum_{j \neq n} |H_{n,m}^H D(\alpha)H_{j,m}|^2 / x_j + \sigma_j^2 / \beta_n \right)} \right)$$

and

$$\pi(p, \alpha) \rightarrow \tilde{\pi}(x, \alpha) \triangleq \sum_{n=1}^{N} \frac{\tilde{h}_n(D(\alpha))^2 \tilde{h}_n}{x_n} = \sum_{n=1}^{N} \sum_{m=1}^{M} \frac{|\tilde{h}_{n,m}|^2}{(\lambda_m + \alpha)^2 x_n}$$

which is seen as a convex function by checking its Hessian.

Then (13) is written by the following convex constrained optimization problem:

$$\max_{x \in \mathbb{R}_+^N, \alpha > 0} \min_{n=1, \ldots, N} \tilde{r}_n(x, \alpha)$$  \quad (16a)

subject to

$$\tilde{\pi}(x, \alpha) \leq P. \quad (16b)$$

The difficulty of computing (16) is now concentrated at its objective function, which is still a very complex nonconcave function. To treat it, introduce the new slack variable $z = (z_1, \ldots, z_M)^T \in \mathbb{R}_+^M$ constrained by

$$z_m = \frac{1}{\lambda_m + \alpha}, m = 1, \ldots, M, \quad (17)$$

which is equivalent to the nonconvex constraint

$$z_m \geq \frac{1}{\lambda_m + \alpha}, m = 1, \ldots, M. \quad (18)$$

plus the convex constraint

$$z_m(\lambda_m + \alpha) \geq 1, m = 1, \ldots, M. \quad (19)$$

As the first stage, we treat the following nonconvex relaxed problem by dropping the nonconvex constraints (18):

$$\max_{(x, z) \in \mathbb{R}_+^{N+M}, \alpha > 0} \Phi(x, \alpha, z) \triangleq \min_{n=1, \ldots, N} g_n(x, \alpha, z)$$  \quad (20a)

subject to

$$z_m(\lambda_m + \alpha) \geq 1, m = 1, \ldots, M. \quad (19)$$

(18b)

where

$$g_n(x, \alpha, z) \triangleq \ln \left( 1 + \frac{\left( \sum_{m=1}^{M} |\tilde{h}_{n,m}|^2 / (\lambda_m + \alpha) \right)^2}{x_n \left( \sum_{j \neq n} |\tilde{h}_{n,m}^* \tilde{h}_{j,m} z_m|^2 / x_j + \sigma_j^2 / \beta_n \right)} \right) \quad (21)$$

Suppose $(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})$ is a feasible point for (20) found from the $(\kappa - 1)$th iteration. Use inequality (78) in the Appendix for

$$x = x_n, y = 1 / \left( \sum_{m=1}^{M} |\tilde{h}_{n,m}|^2 / (\lambda_m + \alpha) \right)^2,$$

$$t = \left( \sum_{j \neq n} \sum_{m=1}^{M} \tilde{h}_{n,m}^* \tilde{h}_{j,m} z_m \right)^2 / x_n + \sigma_j^2 / \beta_n \right)$$

and

$$\bar{x} = x^{(\kappa)}, \bar{y} = y^{(\kappa)} \triangleq 1 / \left( \sum_{m=1}^{M} |\tilde{h}_{n,m}|^2 / (\lambda_m + \alpha) \right)^2,$$

$$\bar{t} = \bar{t}^{(\kappa)} \triangleq \left( \sum_{j \neq n} \sum_{m=1}^{M} \tilde{h}_{n,m}^* \tilde{h}_{j,m} z_m \right)^2 / x_n + \sigma_j^2 / \beta_n \right),$$
to obtain
\[ g_n(x, \alpha, z) \geq g_n^{(\kappa)}(x, \alpha, z) \]
over the trust region
\[ \lambda_m + 2\alpha^{(\kappa)} - \alpha > 0, \quad m = 1, \ldots, M, \]
for
\[ g_n^{(\kappa)}(x, \alpha, z) = \frac{\tilde{\alpha}_n^{(\kappa)} - \tilde{b}_n^{(\kappa)} x_n - \sum_{j \neq n} \tilde{h}_{n,m}^{(\kappa)} z_{m,z}^{2} + \frac{\sigma^2}{\beta_n} \}, \]
which is a concave function with
\[ 0 < \tilde{a}_n^{(\kappa)} \triangleq \ln \left(1 + 1/x_n \right) y_n^{(\kappa)} l_n^{(\kappa)} + 3, \]
\[ 0 < \tilde{b}_n^{(\kappa)} \triangleq 1/x_n y_n^{(\kappa)} l_n^{(\kappa)} + 1, \]
\[ 0 < \tilde{c}_n^{(\kappa)} \triangleq 1/x_n y_n^{(\kappa)} l_n^{(\kappa)} + 1, \]
\[ 0 < \tilde{d}_n^{(\kappa)} \triangleq 1/x_n y_n^{(\kappa)} l_n^{(\kappa)} + 1, \]
\[ \ell_n^{(\kappa)}(\alpha) = (\lambda_m + 2\alpha^{(\kappa)} - \alpha)/\lambda_m + \alpha^{(\kappa)} \cdot 2. \]
At the \( \kappa \)th iteration, the following convex program is solved to generate the next feasible point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\) for (20):
\[
\max_{(x,z) \in \mathbb{R}^N_{+,M}, \alpha > 0} \Phi^{(\kappa)}(x, \alpha, z) \triangleq \min_{n=1, \ldots, N} g_n^{(\kappa)}(x, \alpha, z) \tag{27a}
\]
s.t. \((16b), (19), (23)\). \(27b\)
The computational complexity of (27) is
\[ \mathcal{O}(\tilde{n}^2\tilde{m}^{2.5} + \tilde{m}^{3.5}), \]
with \( \tilde{n} = N + M + 1 \), which is the number of scalar variables, and \( \tilde{m} = 2 + 3M + N \), which is the number of constraints.

Note that \[ \Phi^{(\kappa)}(x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)}) > \Phi^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}) \]
because \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\) and \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\) respectively are the optimal solution and a feasible point for the convex optimization problem (27). This together with \[ \Phi^{(\kappa)}(x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)}) \geq \Phi^{(\kappa)}(x^{(\kappa+1)}, \alpha^{(\kappa)}, z^{(\kappa+1)}), \]
which follows from (22) yield
\[ \Phi^{(\kappa)}(x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)}) \geq \Phi^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}). \]
But one can check immediately that \[ \Phi^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}) \geq \Phi^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}), \]
so
\[ \Phi^{(\kappa)}(x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)}) \geq \Phi^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}), \]
showing that \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\) is a better feasible point for (20) than \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\). The objective \( \Phi^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}) \) thus converges. The sequence \( \{(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\} \) of improved feasible points for (20) can be proved converged at least to a stationary point \((x^{opt}, \alpha^{opt}, z^{opt})\) of (20) \cite{31}.

If \((x^{opt}, \alpha^{opt}, z^{opt})\) satisfies the nonconvex constraint (18) with some tolerance then it is also a stationary point of (16). Otherwise, we go to the second stage, where we follow \cite{24, 32, 33} by introducing the penalty function
\[ \theta_{min}(\alpha, z) = \min_{m=1, \ldots, M} \left[ 1/(\lambda_m + \alpha) - z_m \right], \]
for which \( \theta_{min}(\alpha, z^{opt}) < 0 \), and then the penalty factor
\[ \mu \triangleq \min_{n=1, \ldots, N} g_n(x^{opt}, \alpha^{opt}, z^{opt}) \]
\[ \theta_{min}(\alpha^{opt}, z^{opt}) \]
to consider the following penalized optimization problem
\[ \max_{(x,z) \in \mathbb{R}^N_{+,M}, \alpha > 0} \left[ \min_{n=1, \ldots, N} g_n(x, \alpha, z) + \mu \theta_{min}(\alpha, z) \right] \]
s.t. \((16b), (19), (23)\). \(32\)
Note that the convex constraint (19) implies
\[ z_m \geq 1/(\lambda_m + \alpha), \quad m = 1, \ldots, M, \]
so \( -\theta_{min}(\alpha, z) = \max_{m=1, \ldots, M} (z_m - 1/(\lambda_m + \alpha)) \geq 0 \) and \( -\theta_{min}(\alpha, z) = 0 \) if and only if \((\alpha, z)\) satisfies the nonconvex constraint (18). Therefore \( -\theta_{min}(\alpha, z) \) can be used to measure the degree of satisfaction of the nonconvex constraint (18). Instead of handling the nonconvex constraint (18) we incorporate this degree of its satisfaction into the objective function in (32), which is an exact penalty optimization formulation for (16), i.e. they share the same optimal solution with a finite value of \( \mu \) \cite[Chapter 16]{34}. Initialized from \((x^{(0)}, \alpha^{(0)}, z^{(0)}) = (x^{opt}, \alpha^{opt}, z^{opt})\), at the \( \kappa \)th iteration, the following convex optimization problem is solved to generate the next feasible point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\) for (32):
\[ \max_{(x,z) \in \mathbb{R}^N_{+,M}, \alpha > 0} \left[ \min_{n=1, \ldots, N} g_n(x, \alpha, z) + \mu \theta_{min}(\alpha, z) \right] \]
s.t. \((16b), (19), (23), (33)\)
where
\[ \theta_{min}(\alpha, z) \triangleq \min_{m=1, \ldots, M} \ell_m^{(\kappa)}(\alpha) - z_m \]
which a lower bounding concave approximation of \( \theta_{min}(\alpha, z) \). Its computational complexity is the same as that of (27). Analogously to (29), \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\) is seen as a better feasible point for (32) than \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\). The sequence \( \{(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\} \) at least converges to a stationary point \((x^{opt}, \alpha^{opt}, z^{opt})\) of (32) with
\[ z^{opt} = 1/(\lambda_m + \alpha^{opt}), \quad m = 1, \ldots, M. \]
This computational procedure in the second stage terminates when1
\[ \min_{m=1, \ldots, M} \left[ 1/(\lambda_m + \alpha^{(\kappa)}) - z^{(\kappa)} \right] \geq -\epsilon_{tot} \]
\[ 1 \text{For instance } \epsilon_{tot} = 0.1 \]
so
\[ z_m^{(k)} = 1/(\lambda_m + \alpha^{(k)}) \]
for tolerance \( \epsilon_{tot} \). Algorithm 1 sums up our computational procedure.

**Algorithm 1** Two-stage Max-min Optimization Algorithm

1. **Stage I:** Set \( \kappa = 0 \). Take a feasible point \((x^{(0)}, \alpha^{(0)}, z^{(0)})\) for the convex constraints (16b) and (19).
2. **Repeat until convergence:** Solve the convex optimization problem (27) to generate the next feasible point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\). Set \( \kappa := \kappa + 1 \).
3. Set \((\bar{x}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt}) = (x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\). If it satisfies (18) declare that it is a stationary point of (16) and stop the Algorithm. Otherwise go to Stage II.
4. **Stage II:** Set \( \mu \) by (31). Reset \( \kappa \rightarrow 0 \), and \((\bar{x}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt}) \rightarrow (x^{(0)}, \alpha^{(0)}, z^{(0)})\).
5. **Repeat Until (35):** Solve the convex optimization problem (33) to generate the next feasible point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\); Set \( \kappa := \kappa + 1 \).
6. **Output** \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\) as the stationary point of (16).

### B. EE maximization

Similarly, at the first stage, we address the EE maximization problem (14) via its following nonconvex relaxation

\[
\max_{(x, z) \in \mathbb{R}^{n_M+M}, \alpha > 0} \frac{\sum_{n=1}^{N} g_n(x, \alpha, z)}{|\Pi(x, \alpha)|} \tag{36a}
\]

s.t. (16b), (19), \( g_n(x, \alpha, z) \geq \bar{r}, n = 1, \ldots, N \), \( g_n(x, \alpha, z) \geq \bar{r}, n = 1, \ldots, N \) \tag{36b}

where \( g_n(x, \alpha, z) \) is defined from (21) and \( |\Pi(x, \alpha)| = \eta \tilde{\Pi}(x, \alpha) + P_{non} \).

Suppose that \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\) is a feasible point for (36) found from the \( (\kappa-1) \)th iteration and

\[
\delta^{(\kappa)} \triangleq \frac{\sum_{n=1}^{N} g_n(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})}{|\Pi(x^{(\kappa)}, \alpha^{(\kappa)})|} \tag{37}
\]

At the \( \kappa \)th iteration we solve the following convex optimization problem to generate the next iterative point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\):

\[
\max_{(x, z) \in \mathbb{R}^{n_M+M}, \alpha > 0} \sum_{n=1}^{N} g_n^{(\kappa)}(x, \alpha, z) - \delta^{(\kappa)}|\Pi(x, \alpha)| \tag{38a}
\]

s.t. (16b), (19), (23), \( g_n^{(\kappa)}(x, \alpha, z) \geq \bar{r}, n = 1, \ldots, N \) \tag{38b}

where functions \( g_n^{(\kappa)} \) are derived from (24). Its computational complexity is (28) for \( n = M + N + 1 \) and \( m = 2 + 3M + 2N \). Since \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\) is a feasible point for (38), it follows that

\[
\sum_{n=1}^{N} g_n^{(\kappa)}(x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)}) - \delta^{(\kappa)}|\Pi(x^{(\kappa+1)}, \alpha^{(\kappa+1)})| > \sum_{n=1}^{N} g_n^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}) - \delta^{(\kappa)}|\Pi(x^{(\kappa)}, \alpha^{(\kappa)})| = 0.
\]

Hence

\[
\frac{\sum_{n=1}^{N} g_n^{(\kappa)}(x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})}{|\Pi(x^{(\kappa+1)}, \alpha^{(\kappa+1)})|} > \delta^{(\kappa)},
\]

so

\[
\delta^{(\kappa+1)} = \frac{\sum_{n=1}^{N} g_n(x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})}{|\Pi(x^{(\kappa+1)}, \alpha^{(\kappa+1)})|} \geq \frac{\sum_{n=1}^{N} g_n^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})}{|\Pi(x^{(\kappa)}, \alpha^{(\kappa)})|} > \delta^{(\kappa)},
\]

i.e. the next iterative point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\) is better than \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\) in optimizing (36). Suppose \((\bar{x}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})\) is the stationary point of (36) found by this procedure. Again, if \((\bar{x}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})\) satisfies the convex constraint (18) then it is also a stationary point for the EE maximization problem (14). Otherwise, we initialize the second stage with setting

\[
\mu = -\sum_{n=1}^{N} \frac{g_n(\bar{x}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})/\theta_{sum}(\bar{x}^{opt}, \bar{z}^{opt})}{\theta_{sum}(\bar{x}^{opt}, \bar{z}^{opt})},
\]

where

\[
\theta_{sum}(\alpha, z) = \sum_{m=1}^{M} \left( \frac{1}{\lambda_m + \alpha} - z_m \right). \tag{41}
\]

Initialized from \((x^{(0)}, \alpha^{(0)}, z^{(0)}) = (\bar{x}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})\), at the \( \kappa \)th iteration we solve the following convex optimization problem to generate the next iterative point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\):

\[
\max_{(x, z) \in \mathbb{R}^{n_M+M}, \alpha > 0} \sum_{n=1}^{N} g_n^{(\kappa)}(x, \alpha, z) + \mu \delta^{(\kappa)}|\Pi(x, \alpha)| \tag{42a}
\]

s.t. (16b), (19), (23), \( g_n^{(\kappa)}(x, \alpha, z) \geq \bar{r}, n = 1, \ldots, N \) \tag{42b}

where

\[
\delta^{(\kappa)} = \frac{\sum_{m=1}^{M} (\ell_m^{(\kappa)}(\alpha) - z_m)}{\theta_{sum}(\alpha, z)} \tag{43}
\]

with \( \ell_m^{(\kappa)} \) defined from (26) and

\[
\delta^{(\kappa)} \triangleq \frac{\sum_{n=1}^{N} g_n^{(\kappa)}(x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}) + \mu \theta(\alpha^{(\kappa)}, z^{(\kappa)})}{|\Pi(x^{(\kappa)}, \alpha^{(\kappa)})|}. \tag{44}
\]

Its computational complexity is the same as that of (38). Algorithm 2 sums up our computational procedure.

### III. MIMO REGULARIZED ZERO-FORCING

Now, consider the case when each UE is equipped with \( N_r \) antennas and the number of information streams for each use is \( N_r \) so each information \( s_i \in \mathbb{C}^{N_r}, E[|s_i|^2] = 1 \) for UE \( i \) is precoded by \( P_n \in \mathbb{C}^{M \times N_r} \).

Instead of MISO equation (1), we have the MIMO equation

\[
y = \text{diag}[\sqrt{\beta_n}I_{N_r}] \sum_{n=1}^{N} H^H s + w, \tag{45}
\]
Algorithm 2 Two-stage EE Maximization Algorithm

1: **State I:** Use Algorithm 1 to find a feasible point \((x^{(0)}, \alpha^{(0)}, z^{(0)})\) for constraints (16b), (19) and (36b). Set \(\kappa = 0\).

2: **Repeat until convergence:** Solve the convex optimization problem (38) to generate the next iterative point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\); Set \(\kappa := \kappa + 1\).

3: Set \((\bar{x}^{opt}, \alpha^{opt}, \bar{z}^{opt}) = (x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\). If it satisfies (18) declare that it is a stationary point of (14) and stop the Algorithm. Otherwise go to the next stage.

4: **Stage II:** Set \(\mu\) by (40). Reset \(\kappa \rightarrow 0\), and \((\bar{x}^{opt}, \alpha^{opt}, \bar{z}^{opt}) = (x^{(0)}, \alpha^{(0)}, z^{(0)})\).

5: **Repeat Until (35):** Solve the convex optimization problem (42) to generate the next feasible point \((x^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\); Set \(\kappa := \kappa + 1\).

6: **Output** \((x^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\) as the stationary point of (14).

where

\[
\mathbf{H}^H = \begin{bmatrix} h_{1,1}^*, \ldots, h_{N_N,1}, \ldots, h_{N_N,N_N}^* \end{bmatrix} \in \mathbb{C}^{(N_N,N_N) \times N},
\]

and \(y = (y_1^T, \ldots, y_N^T)^T \in \mathbb{C}^{N \times N}, y_i \in \mathbb{C}^{N_r}, \) is the signal received at \(N\) UEs, \(s = \sum_{i=1}^{N} p_i s_i \in \mathbb{C}^{M}\) is the signal sent from the BS, and \(w \in \mathbb{C}^{N \times (0, \sigma^2 I_{N_N})}\) is the background noise. Compared to the channel (2), which is a vector, the channel from the BS to user \(n\) is a matrix

\[
\sqrt{\beta_n} h_{n},
\]

where \(\sqrt{\beta_n}\) expresses the large-scale fading as in (2), while entries of \(h_{n} = R_{T,n}^H h \circ R_{R,n}\) with \(h_{ij} \in \mathcal{C}(0,1)\) expressing the small scale fading, \(h_{n} = [h_{1,1}, \ldots, h_{1,M}; \ldots; h_{N_N,1}, \ldots, h_{N_N,M}]^T\) and \(R_{T,n}\) and \(R_{R,n}\) expressing the spatial correlation at BS and UE \(n\), respectively.

We consider the following new class of RZF

\[
\hat{P} = \mathbf{H} (\mathbf{H}^H + \alpha \mathbf{I}_{N_N})^{-1} \text{diag}(\mathbf{p}_n)^N, \quad \mathbf{p}_n = \text{diag}(\sqrt{\beta_n} I_{N}),
\]

which is a Hermitian symmetric matrix. Each its block entry is

\[
\mathbf{H}_{n}^H (\mathbf{H}^H + \alpha \mathbf{I}_M)^{-1} \mathbf{H}_{n} \in \mathbb{C}^{N_r \times N_r}.
\]

Upon the SVD (8), define

\[
\bar{h}_{n} = \mathbf{U}^H h_{n}, n = 1, \ldots, N.
\]

Then,

\[
\mathcal{H}_\alpha \triangleq \frac{[\bar{h}_{n}^H D(\alpha) \bar{h}_{n}^j]_{n=1, \ldots, N; j=1, \ldots, N} \in \mathbb{C}^{(N_N,N_N)}}{(\mathcal{D}(\alpha) \bar{h}_{n})},
\]

where \(\mathcal{D}(\alpha)\) is defined from (10).

Then, \(r_n(P, \alpha) = \ln |\mathbf{P}_n^2| + 2 \ln |\mathbf{H}_{n}^H D(\alpha) h_{n}| + \ln \left| (\mathbf{H}_{n}^H D(\alpha) h_{n}^j)^2 + \sigma^2 I_{N_r} \right|^{-1}ceil (49)

while the power constraint at the BS is

\[
\pi(P, \alpha) = \sum_{i=1}^{N} \text{trace} (\mathbf{H}_{n}^H D(\alpha) h_{n}^j)^2 \mathbf{H}_{n}^2 (P, \alpha) \leq P.
\]

The max-min rate optimization is now formulated as

\[
\max_{P > 0, \alpha > 0, n=1, \ldots, N} r_n(P, \alpha) \quad \text{s.t.} \quad \pi(P, \alpha) \leq P.
\]

while the EE maximization problem is

\[
\max_{P > 0, \alpha > 0, n=1, \ldots, N} \sum_{i=1}^{N} r_n(P, \alpha) \quad \text{s.t.} \quad \pi(P, \alpha) \leq P.
\]

Compared to their counterparts for single-antenna UEs defined by (11) and (12), the throughput and sum power functions for multi-antenna UEs defined by (49) and (50) are even more complex. Obviously, the tools for solving (14) and (13) in the previous section are hardly usable for solving (52) and (53). The next subsections provide a new development for their computation.

A. MIMO max-min throughput optimization

There are lots of ways to handle the optimization problem (52). Perhaps, a natural way is to introduce the variables

\[
X_n \triangleq \mathbf{P}_n^2 = \text{diag}[1/p_{n,i}]_{i=1, \ldots, N_r}, \quad \mathbf{X} \triangleq \text{diag}[X_{n}]_{n=1, \ldots, N},
\]

Then

\[
r_n(P, \alpha) \rightarrow \tilde{r}_n(X, \alpha) = \ln |X_n^{-1}| + 2 \ln |\mathbf{H}_{n}^H D(\alpha) h_{n}| + \ln \left| \Phi_{nn}(X, \alpha) + \frac{\sigma^2 I_{N_r}}{\tilde{r}_n} \right|^{-1}ceil (55)
\]
and
\[
\pi(P, \alpha) \rightarrow \bar{\pi}(X, \alpha) = \sum_{n=1}^{N} \text{trace}((H_n^H(D(\alpha))^2H_nX_n^{-1}))
\]
\[
= \sum_{n=1}^{N} \sum_{i=1}^{N_r} \sum_{m=1}^{M} \frac{|H_n(m,i)|^2}{(\lambda_m + \alpha)^2 x_{ni}},
\]
which is a convex function for
\[
\Phi_nj(X_j, \alpha) \triangleq (H_n^HD(\alpha)H_j)X_j^{-1}(H_n^HD(\alpha)H_j)^H,
\]
\[
n = 1, \ldots, N; j = 1, \ldots, N,
\]
and
\[
\Psi_n(X, \alpha) \triangleq \sum_{j \neq n} \Phi_nj(X_j, \alpha).
\]
The max-min rate optimization (52) is now equivalent to
\[
\max_{X > 0, \alpha > 0} \min_{n=1, \ldots, N} \bar{r}_n(X, \alpha)
\]
s.t. \(\bar{\pi}(X, \alpha) \leq P,\)
(59a)
\[
\begin{align*}
\max_{X > 0, z \in \mathbb{R}^M_{+}, \alpha > 0} & \min_{n=1, \ldots, N} g_n(X, \alpha, z) \\
\text{s.t.} & \quad (19), (59b),
\end{align*}
\]
where
\[
g_n(X, \alpha, z) = \ln |X_n^{-1}| + 2 \ln |H_n^H(\Sigma + \alpha I_M)^{-1}H_n| + \\
+ \ln \left( |(H_n^H\Delta(z)H_n)X_n^{-1}(H_n^H\Delta(z)H_n)^H|^{-1} \right) - \\
\left( \sum_{j \neq n} |(H_n^H\Delta(z)H_j)X_j^{-1}(H_n^H\Delta(z)H_j)^H|^{-1} \right) - \\
\left( \sum_{j \neq n} |(H_n^H\Delta(z)H_j)X_j^{-1}(H_n^H\Delta(z)H_j)^H|^{-1} \right) + \frac{\sigma^2}{\beta_n} I_{N_r},
\]
(61)
for \(\Delta(z) \triangleq \text{diag}[z_m^M]_{m=1}^M.\)
Suppose that \((X^{(k)}(\alpha), \alpha^{(k)}, z^{(k)})\) is a feasible point for (60) found from the \((\kappa - 1)\)th iteration.
We address the first term in the objective (61). By applying inequality (82) in the Appendix for \(X_1 = X_n\) and \(X_1 = X^{(k)}\), we obtain
\[
\ln |X_n^{-1}| \geq g_{1,n}^{(k)}(X)
\]
for
\[
g_{1,n}^{(k)}(X) \triangleq \log |(X^{(k)}(\alpha))^{-1}| + \text{trace}(I) - \text{trace}((X^{(k)}(\alpha))^{-1}X_n).
\]
We address the second term in the objective (61). By applying inequality (83) in the Appendix for \(\alpha > 0\) and \(\alpha = \alpha^{(k)} > 0\), we obtain
\[
\ln |H_n^H(\Sigma + \alpha I_M)^{-1}H_n| \geq g_{2,n}^{(k)}(\alpha)
\]
for
\[
g_{2,n}^{(k)}(\alpha) = \ln |H_n^H(\Sigma + \alpha(\kappa) I_M)^{-1}H_n| - \\
\text{trace}\left( (H_n^H(\Sigma + \alpha(\kappa) I_M)^{-1}H_n)^{-1} \times \\
H_n^H(\Sigma + \alpha(\kappa) I_M)^{-2}H_n \right) (\alpha - \alpha^{(k)}).
\]
We address the third term in the objective (61). By applying inequality (81) in the Appendix for
\[
X_1 = (H_n^H\Delta(z)H_n)X_n^{-1}(H_n^H\Delta(z)H_n)^H,
\]
\[
X_2 = \sum_{j \neq n} (H_n^H\Delta(z)H_j)X_j^{-1}(H_n^H\Delta(z)H_j)^H + (\sigma^2/\beta_n)I_{N_r},
\]
and
\[
\tilde{X}_1 = (H_n^H\Delta(z^{(k)})H_n)(X^{(k)}(\alpha))^{-1}(H_n^H\Delta(z^{(k)})H_n)^H,
\]
\[
\tilde{X}_2 = \sum_{j \neq n} (H_n^H\Delta(z^{(k)})H_j)(X^{(k)}(\alpha))^{-1}(H_n^H\Delta(z^{(k)})H_j)^H
\]
we obtain
\[
\ln |X_1^{-1} + X_2^{-1}| \geq g_{3,n}(X, z)
\]
for
\[
g_{3,n}(X, z) = \log |(X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}| + \text{trace}(I) + \\
+ \text{trace}([((X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1})]X_1) + \\
+ \text{trace}([((X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1})]X_2).
\]
For \(A \succeq 0\), it can be derived that
\[
\text{trace}(A(X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}) = \\
\text{trace}(A(X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}) + \\
\text{trace}(A(X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}) A \text{diag}[1/p_{ji}]_{i=1}^{N_r} = \\
\sum_{i=1}^{N_r} ||(H_n^A)X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}||/p_{ji} = \\
\sum_{i=1}^{N_r} ||(H_n^A)X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}||/p_{ji},
\]
where \(H_j = H_j(\cdot, i) \in \mathbb{C}^M\) is the ith column of \(H_j\).
Therefore, for 0 \leq A_i \triangleq (X^{(k)}(\alpha))^{-1} - (X_1^{(k)}(\alpha))^{-1} - (X_2^{(k)}(\alpha))^{-1},
\]
i = 1, 2, one has
\[
g_{3,n}^{(k)}(X, z) = \log |(X^{(k)}(\alpha))^{-1} + (X^{(k)}(\alpha))^{-1}| + \\
\text{trace}(I) - \text{trace}(A_2) \sigma^2/\beta_n + \\
- \sum_{i=1}^{N_r} ||(H_n^A)X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}||/p_{ni} = \\
- \sum_{j \neq n} ||(H_n^A)X_1^{(k)}(\alpha))^{-1} + (X_2^{(k)}(\alpha))^{-1}||/p_{ji},
\]
which is a concave function.
At the $k$th iteration, the following convex program is solved to generate the next feasible point $(X^{(k+1)}, \alpha^{(k+1)}, z^{(k+1)})$ for (59):

$$\max_{X>0, z \in \mathbb{R}^M, \alpha} \min_{n=1, \ldots, N} g_n^{(\kappa)}(X, \alpha, z) \quad s.t. \; (19), (59b), \quad (68)$$

where $g_n^{(\kappa)}(X, \alpha, z) = g_n^{(\kappa)}(X) + 2g_{2,n}^{(\kappa)}(\alpha) + g_{3,n}^{(\kappa)}(X, z)$. Its computational complexity is (28) for $\bar{n} = NN_r + M + 1$ and $\bar{n} = NN_r + M + 2$.

Similar to the first stage of Algorithm 1, we can show that the sequence $\{(X^{(k)}, \alpha^{(k)}, z^{(k)})\}$ generated by a such procedure will at least converge to a stationary point $(X^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})$ of the nonconvex optimization problem (60). If $(X^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})$ satisfies the nonconvex constraint (18) with some tolerance then it is also a stationary point of (59). Otherwise, recalling the definition (30) of the penalty function $\theta_{min}(\alpha, z)$, we define the penalty factor

$$\mu \triangleq \min_{n=1, \ldots, N} \frac{g_n(\bar{X}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})}{\theta_{min}(\bar{\alpha}^{opt}, \bar{z}^{opt})} \quad (69)$$

to consider the following penalized optimization problem

$$\max_{X>0, z \in \mathbb{R}^M, \alpha>0} \left[ \min_{n=1, \ldots, N} g_n(X, \alpha, z) + \mu \theta_{min}(\alpha, z) \right] \quad s.t. \; (19), (59b), \quad (70a)$$

$$\max_{X>0, z \in \mathbb{R}^M, \alpha>0} \left[ \min_{n=1, \ldots, N} g_n(X, \alpha, z) + \mu \theta^{(\kappa)}_{min}(\alpha, z) \right] \quad s.t. \; (19), (59b), \quad (71a)$$

where $\theta^{(\kappa)}(\alpha, z)$ is defined from (34). Its computational complexity is the same as that of (59). The procedure terminates once (35) is verified.

Algorithm 3 sums up our computational procedure.

Algorithm 3 Two-stage MIMO Max-min Optimization Algorithm

1: **Stage I:** Set $\kappa = 0$. Take a feasible point $(X^{(0)}, \alpha^{(0)}, z^{(0)})$ for the convex constraints (19) and (59b).
2: **Repeat until convergence** Solve the convex optimization problem (68) to generate the next feasible point $(X^{(k+1)}, \alpha^{(k+1)}, z^{(k+1)})$; Set $\kappa := k + 1$.
3: Set $(X^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt}) = (X^{(k)}, \alpha^{(k)}, z^{(k)})$. If it satisfies (18) declare that it is a stationary point of (59) and stop the Algorithm. Otherwise go to the next stage.
4: **Stage II:** Set $\mu$ by (69). Reset $\kappa \to 0$, and $(X^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt}) \to (X^{(0)}, \alpha^{(0)}, z^{(0)})$.
5: **Repeat until (35):**
6: Solve the convex optimization problem (71) to generate the next feasible point $(X^{(k+1)}, \alpha^{(k+1)}, z^{(k+1)})$;
7: Set $\kappa := k + 1$.
8: **Output** $(X^{(k)}, \alpha^{(k)}, z^{(k)})$ as a stationary point of (59).

**B. EE maximization**

At the first stage we address the EE maximization problem (53) via its following nonconvex relaxation

$$\max_{X>0, z \in \mathbb{R}^M, \alpha>0} \sum_{n=1}^{N} g_n(X, \alpha, z) \quad \Pi(X, \alpha) \quad (72a)$$

$$\max_{X>0, z \in \mathbb{R}^M, \alpha>0} \sum_{n=1}^{N} g_n(X, \alpha, z) - \delta(n) \Pi(X, \alpha) \quad \Pi(X, \alpha) \quad (74a)$$

s.t. (19), (59b), $g_n(X, \alpha, z) \geq \bar{r}, n = 1, \ldots, N$. (72b)

Its computational complexity is (28) for $\bar{n} = NN_r + M + 1$ and $\bar{n} = 2NN_r + M + 2$. Similar to (39), we can show that $\delta(n) > \delta(n+1)$ and so the generated sequence $\{(X^{(k)}, \alpha^{(k)}, z^{(k)})\}$ at least converges to a stationary point $(X^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})$ of (72). Again, if $(X^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})$ satisfies the convex constraint (18) then it is also a stationary point of the EE maximization problem (53). Otherwise, we need to process one more stage to compute a stationary point of (53). To this end, with $\theta_{sum}$ defined from (41), set

$$\mu = -\sum_{n=1}^{N} g_n(\bar{X}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})/\theta_{sum}(\bar{\alpha}^{opt}, \bar{z}^{opt}). \quad (75)$$

Its computational complexity is the same as that of (74). Algorithm 4 sums up our computational procedure.
Algorithm 4 MIMO EE Maximization Algorithm

1: **Initialization:** Use Algorithm 3 to find a feasible point \((X^{(0)}, \alpha^{(0)}, z^{(0)})\) for constraints (19), (59b), and 
\(g_n(X, \alpha, z) \geq \tilde{r}_n, n = 1, \ldots, N.\) Set \(\kappa = 0.\)

2: **Repeat until convergence:** Solve the convex optimization problem (74) to generate the next iterative point 
\((X^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\). Set \(\kappa := \kappa + 1.\)

3: Set \((\bar{X}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt}) = (X^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}).\) If \((\bar{X}^{opt}, \bar{\alpha}^{opt}, \bar{z}^{opt})\) satisfies (18) declare that it is a stationary point of (53) and stop the Algorithm. Otherwise set \(\mu\) by (75). Reset \(\kappa \rightarrow 0,\) and 
\((X^{(0)}, \alpha^{(0)}, z^{(0)}) = (X^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)}).\)

4: **Repeat until 35:** Solve the convex optimization problem (76) to generate the next feasible point 
\((X^{(\kappa+1)}, \alpha^{(\kappa+1)}, z^{(\kappa+1)})\). Set \(\kappa := \kappa + 1.\)

5: **Output** \((X^{(\kappa)}, \alpha^{(\kappa)}, z^{(\kappa)})\) as a stationary point of (53).

IV. Numerical Simulations

In this section, we evaluate the performance of the proposed algorithms by numerical examples. The BS is located at the centre of a hexagon cell with radius 1 km. Unless otherwise stated, it is assumed that half of number of UEs called by far UEs are equally distributed at farther distances with another half of UEs called by nearer UEs equally distributed at near distances to the BS.

We adopt the standard exponential correlation model [16], where the medium spatial correlation between antenna \((p, q)\) and antenna \((m, n)\), which constitute an entry of matrix \(R\) in (2) is expressed by

\[
[R]_{(p,q),(m,n)} = (0.5)^{|p-m|+|q-n|}
\]

Table I provides other simulation parameters for generating large-scale fading channel and power consumption, which are similar to those used in [35]. The throughput threshold for all UEs is \(\tilde{r} \in \{0.4, 0.8\} \) bps/Hz as in [36, Table I]. In the simulations, Opt. RZF is referred to the RZFB found by the proposed Algorithms 1-4, while Sub. RZF is referred to the RZFB found by the proposed Algorithms 1-4 for the regularization parameter \(\alpha\) fixed at 40dB as proposed in [11], [37]. Also, Equip-power RZF is referred to the RZFB under equal-power allocation, while Conv. RZF is referred to the conventional RZFB defined from (4), which is computed by adjusting Algorithm 1-4. Each point of the numerical results is the average of 1,000 random channel realizations.

### TABLE I: Large scale fading Setup

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Numerical value</th>
</tr>
</thead>
<tbody>
<tr>
<td>Carrier frequency / Bandwidth</td>
<td>2GHz / 10MHz</td>
</tr>
<tr>
<td>BS transmission power</td>
<td>46 dBm</td>
</tr>
<tr>
<td>Path loss from BS to UE</td>
<td>128.1 + 37.6log_{10}R [dB], R in km</td>
</tr>
<tr>
<td>Shadowing standard deviation</td>
<td>8 dB</td>
</tr>
<tr>
<td>Noise power density</td>
<td>-174 dBm/Hz</td>
</tr>
<tr>
<td>Noise figure</td>
<td>9 dB</td>
</tr>
<tr>
<td>Drain efficiency of amplifier</td>
<td>(\alpha = 1/0.388)</td>
</tr>
<tr>
<td>Circuit power per antenna</td>
<td>(P_A = 189) mW</td>
</tr>
<tr>
<td>Non-transmission power</td>
<td>(P_{C} = 40) dBm</td>
</tr>
</tbody>
</table>

A. Single-antenna UEs

In this scenario, the BS is equipped with an 4 × 4 (4 rows in the horizontal dimension and 4 columns in the vertical dimension) or 8 × 8 (8 rows in the horizontal dimension and 8 columns in the vertical dimension) uniform planar array (UPA) of antennas. Thus, the total number of antennas at BS is \(M = 16\) or \(M = 64\), respectively.

1) Max-min throughput performance: Table II provides the average numbers of iterations of Algorithm 1 for computing Opt. RZF and Sub. RZF under different numbers of UEs to plot Fig. 1a. Algorithm 1 needs a few iterations, where the second stage for penalized optimization is either not necessary or takes only one iteration. Fig. 1a and Fig. 1b plot the achievable UEs’ max-min throughput by different RZFB schemes versus the number of UEs for \(M = 16\) and \(M = 64\), respectively. In both figures, Opt. RZF obviously outperforms Sub. RZF, showing the importance of optimizing the regularization parameter \(\alpha\). With the number of BS antennas increased to \(M = 64\), the number of served UEs also increases. For example, 16-antenna array BS can offer the QoS of 0.8 bps/Hz to 12 and 16 UEs according to Fig. 1a but 64-antenna array BS offer the same QoS to 54 and 70 UEs according to Fig. 1b. As expected, Equip-power RZF achieves the worst UEs’ max-min throughput. Conv. RZF performs much worse than Sub. RZF and Opt. RZF, which means that regularizing the channel \(H^H\) of small-scale fading in (3) is a much more efficient way than regularizing the channel \(H^H\) incorporating both large-scale and small-scale fadings in (4).

<table>
<thead>
<tr>
<th># of UEs ((N))</th>
<th>16</th>
<th>24</th>
<th>32</th>
<th>40</th>
<th>48</th>
<th>56</th>
<th>64</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub. RZF</td>
<td>5.5</td>
<td>5.1</td>
<td>5.0</td>
<td>4.9</td>
<td>4.8</td>
<td>4.7</td>
<td>4.6</td>
</tr>
<tr>
<td>Opt. RZF</td>
<td>5.5</td>
<td>5.1</td>
<td>5.0</td>
<td>4.9</td>
<td>4.8</td>
<td>4.7</td>
<td>4.6</td>
</tr>
</tbody>
</table>

TABLE II: The average number of iterations for \(M = 16\).
Fig. 1: The achievable users’ max-min throughput for $M = \{16, 64\}$.

2) EE performance: Fig. 2a plots the achievable EE by Opt. RZF and Sub. RZF at QoS threshold $\bar{r} = 0.4$ bps/Hz. As expected, Opt. RZF achieves much better EE compared with Sub. RZF. The gap is very wide for the number of UEs exceeding 12.

Table III: The average iterations of Algorithm 2 to plot Fig. 2b for $M = 16$ and $\bar{r} = 0.8$ bps/Hz.

<table>
<thead>
<tr>
<th># of UEs ($N$)</th>
<th>8</th>
<th>10</th>
<th>12</th>
<th>14</th>
<th>16</th>
<th>18</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub. RZF</td>
<td>8.7</td>
<td>8.9</td>
<td>8.6</td>
<td>10</td>
<td>13</td>
<td>15.4</td>
</tr>
<tr>
<td>Opt. RZF</td>
<td>16.4</td>
<td>17.4</td>
<td>17.6</td>
<td>18.8</td>
<td>17</td>
<td>15.4</td>
</tr>
</tbody>
</table>

Table III shows that the average numbers of iterations of Algorithm 2 for computing Opt. RZF are higher compared to that for computing Sub. RZF. Fig. 2b reveals that Opt. RZF achieves a much better EE than Sub. RZF for a higher QoS threshold $\bar{r} = 0.8$ bps/Hz. Especially, Opt. RZF guarantees this QoS for a higher number of UEs than Sub. RZF does. Opt. RZF is able to serve up to 18 UEs but Sub. RZF cannot serve more than 12 UEs. Fig. 3 justifies this intuition. Sub. RZF consumes almost all the allowable transmit power to serve 12 UEs, while Opt. RZF controls the power much better in serving more UEs.

Fig. 2: EE performance for $M = 16$ and $\bar{r} = \{0.4, 0.8\}$ bps/Hz.

Fig. 3: Transmit power for $M = 16$ and $\bar{r} = 0.8$ bps/Hz.

Table IV: The average iterations of Algorithm 3 to plot Fig. 4 for $M = 64$ and $\bar{r} = 0.8$ bps/Hz.

<table>
<thead>
<tr>
<th># of UEs ($N$)</th>
<th>30</th>
<th>36</th>
<th>42</th>
<th>48</th>
<th>54</th>
<th>60</th>
<th>66</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub. RZF</td>
<td>20</td>
<td>30</td>
<td>40</td>
<td>50</td>
<td>60</td>
<td>70</td>
<td>80</td>
</tr>
<tr>
<td>Opt. RZF</td>
<td>20.5</td>
<td>21</td>
<td>21.5</td>
<td>22</td>
<td>22</td>
<td>22</td>
<td>19.5</td>
</tr>
</tbody>
</table>
For a higher number of BS antennas such as $M = 64$, Fig. 4a demonstrates the superior performance of Opt. RZF compared with Sub. RZF. The average iterations of Algorithm 2 for their computation are given in Table IV. As expected, Opt. RZF serves a much higher number of UEs (up to 66 UEs) than Sub. RZF does at the QoS threshold of $\bar{r} = 0.8$ bps/Hz. This can be explained by the result of transmit power as shown in Fig. 4b, where Opt. RZF requires much less transmission power to optimize EE than Sub. RZF.

### Multi-antenna users

In this scenario, the BS is equipped with an $4 \times 4$ uniform planar array (UPA) of antennas corresponding to $M = 16$ and each UE is equipped with 2 antennas. To see the trade-off between the computational complexity and performance, we also include the simulation for unstructured MIMO beamforming, which uses the algorithm of [25], [26] for the max-min throughput optimization and the algorithm of [38], [39] for the EE maximization. The computational complexity of each iteration is $O(\tilde{n}^2\tilde{m}^{2.5} + \tilde{m}^{3.5})$, where $\tilde{n} = MNr + 1$ and $\tilde{m} = N + 1$ for max-min throughput optimization and $\tilde{n} = N(M^2 + N^2)$ and $\tilde{m} = 2N + M^2N + 1$ for the EE optimization, which is much more complex than that of each iteration of Algorithms 3 and 4.

Table V provides the average numbers of iterations for plotting Fig. 5. Although Opt. RZF requires a larger number of iterations than Sub. RZF for convergence, the worst UEs’ throughput of the former is much higher compared to the latter as shown in Fig. 5. The performance of the unstructured beamforming is slightly higher than that by Opt. RZF.

Fig. 5 plots the achievable UEs’ max-min throughput by different RZFB schemes vs the number of UEs. Opt. RZF is seen achieving the UEs’ max-min throughput much better than Sub. RZF, especially for $N < 20$. Like Fig. 1a and Fig. 1b, the achievable UE’s max-min throughput by Equip-power RZF is the worst and Conv. RZF performs much worse than Sub. RZF and Opt. RZF.

### TABLE V: The average iterations of Algorithm 3 and Unstruct. Max-Min for plotting Fig. 5.

<table>
<thead>
<tr>
<th># of UEs ($N$)</th>
<th>8</th>
<th>12</th>
<th>16</th>
<th>20</th>
<th>24</th>
<th>28</th>
<th>32</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub. RZF</td>
<td>25.3</td>
<td>36.5</td>
<td>45.2</td>
<td>38.6</td>
<td>36</td>
<td>21.8</td>
<td>15.3</td>
</tr>
<tr>
<td>Opt. RZF</td>
<td>38.8</td>
<td>47.3</td>
<td>48.8</td>
<td>45.4</td>
<td>45.9</td>
<td>35.4</td>
<td>26.3</td>
</tr>
<tr>
<td>Unstruct. Max-Min</td>
<td>25</td>
<td>26</td>
<td>35</td>
<td>41</td>
<td>39</td>
<td>35</td>
<td>22</td>
</tr>
</tbody>
</table>

Fig. 5: The achievable UEs’ max-min throughput.

In terms of EE performance, Fig. 6a and Fig. 6b also show the superior EE of Opt. RZF over Sub. RZF particularly because according to Fig. 7, Opt. RZF requires a much lower transmit power than Sub RZF in guaranteeing QoS threshold $\bar{r} = 0.8$ bps/Hz. The average iterations of Algorithm 4 for their computation are provided in Table VI.
V. Conclusions

In this paper, we have considered the joint design of power allocation and regularization parameter to maximize the UEs’ minimal throughput or the energy-efficiency in guaranteeing the UEs’ QoS in term of their throughput requirement for a new class of RZF. We have developed the two-stage algorithms to solve the posed problems and provided numerical examples to support their efficiency. We have also compared the performance achieved by the propose joint design with that achieved by the existing RZF schemes to show the importance of the joint optimization in power allocation and regularization parameter. An extension of the joint design to multi-cell RZF is under our current study.

APPENDIX I: FUNDAMENTAL INEQUALITIES

Theorem 1: Function \( f(x, y, t) = \ln(1 + 1/xyt) \) is convex in \((x, y, t)^T \in \mathbb{R}^3\).

Proof: Function \( \ln(1 + 1/x^3) \) is convex in \( t > 0 \) as it is immediate to see that \( d^2 \ln(1 + 1/t^3)/dt^2 > 0 \), while function \( g(x, y, z) = (xyt)^{1/3} \) is concave in \((x, y, t)^T \in \mathbb{R}^3\) [40]. Then for \( x \geq 0, y \geq 0, t > 0 \) and \( \alpha + \beta = 1 \), it is true that
\[
\begin{align*}
&\alpha f(x, y, t) + \beta f(x_1, y_1, t_1) \\
\leq &\alpha f(x, y, t) + \beta f(x_1, y_1, t_1) + \alpha \ln(1 + 1/(g(x, y, t)^3)) + \beta \ln(1 + 1/(g(x_1, y_1, t_1)^3))
\end{align*}
\]
showing the convexity of \( f(x, y, t) \).

The following linear approximation then holds true for all \( x > 0, \bar{x} > 0, y > 0, \bar{y} > 0, t > 0 \) and \( \bar{t} > 0 \) [40]:
\[
\ln(1 + 1/xyt) \geq f(\bar{x}, \bar{y}, \bar{t}) + (\nabla f(\bar{x}, \bar{y}, \bar{t}), (x, y, t) - (\bar{x}, \bar{y}, \bar{t}))
\]
\[
= \bar{a} - \bar{b}x - \bar{c}y - \bar{d}t,
\]
where \( \nabla \) is the gradient operation and
\[
\bar{a} = \ln(1 + 1/\bar{x}\bar{y}\bar{t}) + \frac{3}{(\bar{x}\bar{y}\bar{t} + 1)} > 0, \bar{b} = \frac{1}{(\bar{x}\bar{y}\bar{t} + 1)\bar{x}} > 0,
\]
\[
\bar{c} = \frac{1}{(\bar{x}\bar{y}\bar{t} + 1)\bar{y}} > 0, \bar{d} = \frac{1}{(\bar{x}\bar{y}\bar{t} + 1)\bar{t}} > 0.
\]
Furthermore, as function \( x^2/t \) is convex on \( x \) and \( t > 0 \), the following inequality holds true
\[
\frac{x^2}{t} \geq \frac{2\bar{x}}{\bar{t}} x - \frac{\bar{x}^2}{\bar{t}^2} t.
\]

Finally, noting that function \( f(\alpha, t) = h^H(\Sigma + \alpha\mathbf{I}_M)^{-2}h/t \) is convex in \( \alpha > 0 \) and \( t > 0 \), we have the following inequality
\[
\frac{h^H(\Sigma + \alpha\mathbf{I}_M)^{-2}h}{t} \geq \frac{h^H(\Sigma + \bar{\alpha}\mathbf{I}_M)^{-2}h}{\bar{t}} + (\nabla f(\bar{\alpha}, \bar{t}), (\alpha, t) - (\bar{\alpha}, \bar{t}))
\]
\[
= \frac{4}{t} h^H(\Sigma + \alpha\mathbf{I}_M)^{-2}h - \frac{2}{t^2} h^H(\Sigma + \alpha\mathbf{I}_M)^{-3}h - \frac{2}{t^2} h^H(\Sigma + \alpha\mathbf{I}_M)^{-3}h
\]
for all \( \alpha > 0, t > 0, \bar{\alpha} > 0 \) and \( \bar{t} > 0 \).

---

TABLE VI: The average iterations of Algorithm 4 and Unstruct. GEE for plotting Fig. 6b for \( \bar{r} = 0.8 \) bps/Hz.

<table>
<thead>
<tr>
<th># of UEs (N)</th>
<th>4</th>
<th>6</th>
<th>8</th>
<th>10</th>
<th>12</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub. RZF</td>
<td>21.3</td>
<td>26.7</td>
<td>33</td>
<td>39.3</td>
<td>51</td>
</tr>
<tr>
<td>Opt. RZF</td>
<td>34</td>
<td>44.7</td>
<td>43.7</td>
<td>49.3</td>
<td>51</td>
</tr>
<tr>
<td>Unstruct. EE</td>
<td>8</td>
<td>11</td>
<td>12</td>
<td>18</td>
<td>18</td>
</tr>
</tbody>
</table>
APPENDIX II: CONVEXITY OF NONSPECTRAL FUNCTIONS AND LOG-DET INEQUALITIES

In the previous appendix we obtained some results on convexity of function of real vector variables. There are lots of interesting results on convexity of the so-called spectral functions of matrix variables (see e.g. [41]) such as $f(\lambda_{\text{min}}(X))$ or $f(\lambda_{\text{max}}(X))$, where $\lambda_{\text{min}}(X)$ and $\lambda_{\text{max}}(X)$ are respectively the minimal and maximal eigenvalue of the matrix variable $X$, which are dependent on the set of eigenvalues of the matrix variable but not dependent on the order of these eigenvalues. As expected, it is much more difficult to analyse convexity of nonspectral functions and in this Appendix we provide some first results about it.

Theorem 2: For $A > 0$ function $\log |A +HX^{-1}H^H|$ is convex in $X > 0$.

Proof: One has 

$$ (A +HX^{-1}H^H)^{-1} = A^{-1} - A^{-1}(H^HA^{-1}H + X)^{-1}A^{-1} $$

By [42, Appendix B] 

$$ f(X) = A^{-1} - A^{-1}(H^HA^{-1}H + X)^{-1}A^{-1} $$

is concave 

$$ f(\alpha X + \beta Y) > \alpha f(X) + \beta f(Y) $$

for all $X > 0, Y > 0$ and $\alpha 0, \beta > 0, \alpha + \beta = 1$. Therefore 

$$ \log |f(\alpha X + \beta Y)| \geq \log |f(\alpha X)| + \beta \log |f(Y)|, $$

showing that $\log |A +HX^{-1}H^H|^{-1} = -\log |A +HX^{-1}H^H|$ is concave in $X$. \Box

Theorem 3: Function $\log |X_1^{-1} + X_2^{-1}|$ is convex in $X_1 > 0$ and $X_2 > 0$.

Proof: Define 

$$ f_1(X_1, X_2) \triangleq (X_1^{-1} + X_2^{-1})^{-1} = X_1 - X_1(X_1 + X_2)^{-1}X_1. $$

By [42, Appendix B] $f_1$ is concave: 

$$ f_1(\alpha X_1, \beta Y_1 + Y_2) \geq \alpha f_1(X_1, X_2) + \beta f_1(Y_1, Y_2) $$

for all $X_i > 0, Y_i > 0$ and $\alpha 0, \beta > 0, \alpha + \beta = 1$. Therefore 

$$ \log |f_1(\alpha X_1, \beta Y_1 + Y_2)| \geq \log |f_1(X_1, X_2)| + \beta \log |f_1(Y_1, Y_2)|, $$

showing that $\log |f_1(X_1, X_2)| = -\log |X_1^{-1} + X_2^{-1}|$ is concave. \Box

Applying Theorem 3 yields the following inequality 

$$ \log |X_1^{-1} + X_2^{-1}| \geq \log |X_1^{-1} + X_2^{-1}| + \text{trace}(I) + \text{trace}((X_1^{-1} + X_2^{-1})^{-1}X_1^{-1}X_1) + \text{trace}((X_1^{-1} + X_2^{-1})^{-1}X_2^{-1}X_2) $$

for all $X_i > 0, \bar{X}_i > 0, i = 1, 2$. \hspace{1cm} (81)

Using the convexity of function $\log |X^{-1}|$ yields the following inequality 

$$ \log |X_1^{-1}| \geq \log |X_1^{-1}| + \text{trace}(I) - \text{trace}(X_1^{-1}X_1) \quad \forall X_1 > 0, \bar{X}_1 > 0. \hspace{1cm} (82) $$

Noting that function $f(\alpha) = \ln|H_n^H(\Sigma + \alpha I_M)^{-1}H_n|$ is convex in $\alpha > 0$ by Theorem 2, we then have the following inequality 

$$ \ln|H_n^H(\Sigma + \alpha I_M)^{-1}H_n| \geq f(\tilde{\alpha}) + \langle \nabla f(\tilde{\alpha}), \alpha - \tilde{\alpha} \rangle = \ln|H_n^H(\Sigma + \tilde{\alpha} I_M)^{-1}H_n| $$

$$ -\text{trace} \left( H_n^H(\Sigma + \tilde{\alpha} I_M)^{-1}H_n \right)^{-1} \times H_n^H(\Sigma + \alpha I_M)^{-2}H_n (\alpha - \tilde{\alpha}) \hspace{1cm} (83) $$

for all $\alpha > 0$ and $\tilde{\alpha} > 0$.

APPENDIX III: DERIVATION FOR (12)

$$ \pi(\rho, \alpha) = \text{trace}(FH) = \text{trace} \left( \text{diag}[\sqrt{p_{n}}]_{n=1}^{N} (H^H H + \alpha I_N)^{-1} H^H \right) $$

$$ \times H_n^H(\Sigma + \alpha I_M)^{-1} \text{diag}[p_{n}][N] $$

$$ = \text{trace} \left( H_n^H(\Sigma + \alpha I_M)^{-1} \bar{H}_n \right) $$

$$ \times H_n^H(\Sigma + \alpha I_M)^{-2} \bar{h}_n $$

$$ = \sum_{n=1}^{N} p_n (H_n^H + \alpha I_M)^{-2} \bar{h}_n $$

$$ = (12). $$

APPENDIX IV: DERIVATION FOR (49) AND (50)

$$ r_n(P, \alpha) = \ln |I_{N_{r}} + (H_n^H D(\alpha)H_n)^2 + \left( \sum_{j \neq n} (H_n^H D(\alpha)H_j^H P_j) + (\sigma^2 / \beta) I_{N_{r}} \right)^{-1} | = $$

$$ \ln \left( H_n^H D(\alpha)H_n P_n \right)^2 + \ln \left( H_n^H D(\alpha)H_n P_n \right)^{-2} $$

$$ + \left( \sum_{j \neq n} (H_n^H D(\alpha)H_j^H P_j) + (\sigma^2 / \beta) I_{N_{r}} \right)^{-1} $$

$$ = (49), $$

and 

$$ \pi(P, \alpha) = \text{trace}(\hat{P}^H \hat{P}) = \text{trace} \left( \text{diag}[\sqrt{p_{n}}]_{n=1}^{N} (H^H H + \alpha I_{N_{N, r}})^{-1} H^H \right) $$

$$ \times H_n^H(\Sigma + \alpha I_M)^{-1} \text{diag}[p_{n}][N_{n}] $$

$$ = \text{trace} \left( H_n^H(\Sigma + \alpha I_M)^{-1} \bar{H}_n \right) $$

$$ \times (H_n^H + \alpha I_M)^{-1} \text{diag}[p_{n}][N_{n}] $$

$$ = (50). $$

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L. D. Nguyen, H. D. Tuan, T. Q. Duong, O. A. Dobre, and H. V. Poor,
Z. Xu, C. Yang, G. Y. Li, Y. Liu, and S. Xu, “Energy-efficient CoMP
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